PUZZLES WITH SEVERAL CRITICAL POINTS

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Branner-Hubbard Conjecture

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Koslovski-vanStrien and Qiu-Yin proved this conjecture.

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Dynamics

If $f: U \to V$ is a ramified covering of degree d and $A \subset U$, $B \subset V$ are closed disks, $mod(U \setminus A) \ge \frac{1}{d}mod(V \setminus B)$

Property of bounded degree

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Definition

The point x has property (\star) if :

 $\exists z, (k_n)_{n\geq 0} \mid f^{k_n} : P_{k_n+k_0}(x) \rightarrow P_{k_0}(z), \ \forall n \geq 1, \text{ has bounded degree.}$

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Lemma

If a point x has property (\star) then Imp(x) = x.

Definition

For a graph Γ , say that z accumulates y if for every $n \ge 0$ there exists k > 0 such that $f^k(z) \in P_n(y)$. Write $y \in \omega(z)$.

Lemma

Every point x falls in one of the cases:

• If $\omega(x) \cap Crit = \emptyset$ then (*) is satisfied;

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- If ∃c ∈ ω(x) ∩ Crit such that ∃c' ∈ ω(c) ∩ Crit such that c ∉ ω(c') then (*) is satisfied;
- For $\forall c \in \omega(x) \cap Crit$ and $\forall c' \in \omega(c) \cap Crit$ we have that $c \notin \omega(c')$ we call it the recurrent case.

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Definition

 $P_n(c)$ is a successor of a puzzle piece $P_i(c)$ if : $f^{n-i}: P_n(c) = P_i(c)$ and every critical point appears at most twice in $\{f^j(P_n(c)) \mid 0 \le j < n-i\}.$

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Of course if z is non-persistently recurrent then (\star) is satisfied

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Theorem :

There exists a sequence of puzzle pieces (K_n) in the nest $(P_j(c))$ defined by the operator Γ and two operators \mathcal{A} and \mathcal{B} of bounded degree: $K_n := \mathcal{A}(\Gamma^{b+1}(K_{n-1}) \text{ and } K'_n := \mathcal{B}(\Gamma^{b+1}(K_{n-1}) \text{ with the property that} K'_n \setminus K_n \text{ does not intersect the postcritical set.}$

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It is called the enhanced nest by Kozlovski, Shen, van Strien

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• $f^{p_n}(K_n) = K_{n-1}, \ p_{n+1} \ge 2p_n, \ deg(f^{p_n}: K_n \to K_{n-1}) \le C(b, \delta).$

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Using "Kahn-Lyubich covering Lemma", we can prove that $\liminf \mu_n$ is bounded from below.

Hence $Imp(c) = \{c\}.\Box$

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Take some *M*-iterate of them : $f^M : U = K_m \rightarrow V = K_{m-Z}$ (*Z* will be chosen later) and $f^M(y) = z$.

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Pull back by f' the pieces $K_m \subset K'_m$ to $B \subset B'$ containing z, and then by f^M to puzzle pieces $A \subset A'$ containing y.

Claim

- $f^{\xi}(K_{m+2}) \subset A$
- the degree of $f_{|K'_{m+2}}^{\xi}$ is bounded by C_1 independently of Z, m;
- the degree of $f_{|A'|}^M$ is bounded by C_2 independently of Z, m;
- the degree of $f_{|U|}^M$ is bounded by C independent of m;
- the degree of $f'_{|B'}$ is bounded by C_3 independently of Z, m.

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- the degree of $f_{|U|}^M$ is bounded by C independent of m;
- the degree of $f_{|B'|}^{l}$ is bounded by C_3 independently of Z, m.

Hence

•
$$mod(U \setminus A) \leq mod(U \setminus f^{\xi}(K_{m+2})) \leq \frac{1}{C_1}\mu_{m+2}.$$

- $mod(B' \setminus B) \geq \frac{1}{C_3}\mu_m$
- Choose *m* so that $\mu_{m+2} \leq \mu_k$ for $k \leq m+2$.

Hence we get the condition of the covering Lemma:

$$mod(B'\setminus B)\geq rac{1}{C_1C_3}mod(U\setminus A)$$

So that either $mod(U \setminus A) > \epsilon(C_1C_3, D)$ or

 $mod(U \setminus A) \geq \frac{1}{C_1 C_3 2(C_2)^2} mod(V \setminus B).$

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Looking at the first entrance of z in the annuli $K'_i \setminus \overline{K}_i$ for $m \le i \le m - Z + 1$:

 $mod(V \setminus B) \geq \mu_m + \mu_{m-1} + \ldots + \mu_{m-Z+1} \geq Z\mu_{m+2}.$

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So that the first inequality implies that

$$\frac{1}{C_1}\mu_{m+2} \geq \frac{Z}{C_1C_32(C_2)^2}\mu_{m+2}.$$

which is not possible for large Z.

The Covering Lemma

Theorem

Let $f : U \to V$ be a degree D ramified covering. For any $\eta > 0$, there exists $\epsilon = \epsilon(\eta, D) > 0$ such that :

- if $A \subset A' \subset U$ and $B \subset B' \subset V$ are sequences of disks;
- if f is a proper map from A to B, and from A' to B' with degree d;
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Then

- $mod(U \setminus A) > \epsilon$
- or $mod(U \setminus A) > \frac{\eta}{2d^2} mod(V \setminus B)$.

▶ retour

Given a puzzle piece I containing c there exist puzzle pieces A(I) and B(I) containing c such that

- they are pullback of *I*;
- $\mathcal{A}(I) \subset \mathcal{B}(I);$
- the degrees A(I) → I and B(I) → I are bounded by C(b, δ) and one meets c at most b + 1,resp. b times;
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Then define $K'_{n+1} = \mathcal{B}(I_{n+1})$. \frown retour