Bandits for Exploration: Best Arm Identification and Discovery with Probabilistic Experts

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Roadmap

1 Classical Bandits

- 2 Best arm identification in two-armed bandits
 Lower bounds on the complexities
 The complexity of A/B Testing with Gaussian feedback
 The complexity of A/B Testing with binary feedback
- Optimal Exploration with Probabilistic Expert Advice
 Missing mass and Good-UCB
 - Analysis: Classical and Macroscopic Optimality

The (stochastic) Multi-Armed Bandit Model

Environment *K* arms with parameters $\theta = (\theta_1, \dots, \theta_K)$ such that for any possible choice of arm $a_t \in \{1, \dots, K\}$ at time *t*, one receives the reward

 $X_t = X_{a_t,t}$

where, for any $1 \le a \le K$ and $s \ge 1$, $X_{a,s} \sim \nu_a$, and the $(X_{a,s})_{a,s}$ are independent.

Reward distributions $\nu_a \in \mathcal{F}_a$ parametric family, or not. Examples: canonical exponential family, general bounded rewards

Example Bernoulli rewards: $\theta \in [0, 1]^K$, $\nu_a = \mathcal{B}(\theta_a)$

Strategy The agent's actions follow a dynamical strategy $\pi = (\pi_1, \pi_2, ...)$ such that

$$A_t = \pi_t(X_1,\ldots,X_{t-1})$$

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Real challenges

- Randomized clinical trials
 - original motivation since the 1930's
 - dynamic strategies can save resources
- Recommender systems:
 - advertisement
 - website optimization
 - news, blog posts, ...



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- Computer experiments
 - large systems can be simulated in order to optimize some criterion over a set of parameters
 - but the simulation cost may be high, so that only few choices are possible for the parameters
 - Games and planning (tree-structured options)

Classical Bandits

Performance Evaluation, Regret

Cumulated Reward $S_T = \sum_{t=1}^T X_t$

Our goal Choose π so as to maximize

$$\mathbb{E}[S_T] = \sum_{t=1}^T \sum_{a=1}^K \mathbb{E}\left[\mathbb{E}\left[X_t \mathbb{1}\{A_t = a\} | X_1, \dots, X_{t-1}\right]\right]$$
$$= \sum_{a=1}^K \mu_a \mathbb{E}\left[N_a^{\pi}(T)\right]$$

where $N_a^{\pi}(T) = \sum_{t \leq T} \mathbb{1}\{A_t = a\}$ is the number of draws of arm a up to time T, and $\mu_a = E(\nu_a)$.

Regret Minimization equivalent to minimizing

$$R_T = T\mu^* - \mathbb{E}\left[S_T\right] = \sum_{a:\mu_a < \mu^*} (\mu^* - \mu_a) \mathbb{E}\left[N_a^{\pi}(T)\right]$$

where $\mu^* \in \max\{\mu_a : 1 \le a \le K\}$

Asymptotically Optimal Strategies

• A strategy π is said to be consistent if, for any $(\nu_a)_a \in \mathcal{F}^K$,

$$\frac{1}{T}\mathbb{E}[S_T] \to \mu^*$$

The strategy is efficient if for all $\theta \in [0,1]^K$ and all $\alpha > 0$,

$$R_T = o(T^{\alpha})$$

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There are efficient strategies and we consider the best achievable asymptotic performance among efficient strategies

The Bound of Lai and Robbins

One-parameter reward distribution ν_a = $\nu_{\theta_a}, \theta_a \in \Theta \subset \mathbb{R}$.

Theorem [Lai and Robbins, '85]

If π is an efficient strategy, then, for any $\theta \in \Theta^K$,

$$\liminf_{T \to \infty} \frac{R_T}{\log(T)} \ge \sum_{a:\mu_a \le \mu^*} \frac{\mu^* - \mu_a}{\mathrm{KL}(\nu_a, \nu^*)}$$

where $\mathrm{KL}(\nu,\nu')$ denotes the Kullback-Leibler divergence

For example, in the Bernoulli case:

$$KL(\tilde{B}(p),\tilde{B}(q)) = d_{\text{BER}}(p,q) = p\log\frac{p}{q} + (1-p)\log\frac{1-p}{1-q}$$

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Classical Bandits

The Bound of Burnetas and Katehakis

More general reward distributions $\nu_a \in \mathcal{F}_a$

Theorem [Burnetas and Katehakis, '96]

If π is an efficient strategy, then, for any $\theta \in [0,1]^K$,

$$\liminf_{T \to \infty} \frac{R_T}{\log(T)} \ge \sum_{a:\mu_a < \mu^*} \frac{\mu^* - \mu_a}{K_{inf}(\nu_a, \mu^*)}$$



Upper Confidence Bound Strategies

UCB [Lai&Robins '85; Agrawal '95; Auer&al '02]

Construct an upper confidence bound for the expected reward of each arm:



Choose the arm with the highest UCB

- It is an index strategy [Gittins '79]
- Its behavior is easily interpretable and intuitively appealing

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Performance of UCB

For rewards in [0,1], the regret of UCB is upper-bounded as

 $E[R_T] = O(\log(T))$

(finite-time regret bound) and

$$\limsup_{T \to \infty} \frac{\mathbb{E}[R_T]}{\log(T)} \le \sum_{a:\mu_a < \mu^*} \frac{1}{2(\mu^* - \mu_a)}$$

Yet, in the case of Bernoulli variables, the rhs. is greater than suggested by the bound by Lai & Robbins

Many variants have been suggested to incorporate an estimate of the variance in the exploration bonus (e.g., [Audibert&al '07])

The KL-UCB algorithm

The KL-UCB Algorithm, Annals of Statistics 2013 joint work with O. Cappé, O-A. Maillard, R. Munos, G. Stoltz

Parameters: An operator $\Pi_{\mathcal{F}} : \mathfrak{M}_1(\mathcal{S}) \to \mathcal{F}$; a non-decreasing function $f : \mathbb{N} \to \mathbb{R}$

Initialization: Pull each arm of $\{1, \ldots, K\}$ once

for
$$t = K$$
 to $T - 1$ do

compute for each arm a the quantity

$$U_a(t) = \sup \left\{ E(\nu) : \quad \nu \in \mathcal{F} \quad \text{and} \quad KL\left(\Pi_{\mathcal{F}}(\hat{\nu}_a(t)), \nu\right) \leq \frac{f(t)}{N_a(t)} \right\}$$

pick an arm $A_{t+1} \in \underset{a \in \{1, \dots, K\}}{\operatorname{arg max}} U_a(t)$

end for

Exponential Family Rewards

Assume that \(\mathcal{F}_a = \mathcal{F} = \mathcal{canonical exponential family, i.e. such that the pdf of the rewards is given by

$$p_{\theta_a}(x) = \exp\left(x\theta_a - b(\theta_a) + c(x)\right), \quad 1 \le a \le K$$

for a parameter $\theta \in \mathbb{R}^{K}$, expectation $\mu_{a} = \dot{b}(\theta_{a})$

$$U_a(t) = \sup \left\{ \mu \in \overline{I} : \quad d(\hat{\mu}_a(t), \mu) \leq \frac{f(t)}{N_a(t)} \right\}$$

For instance,

for Bernoulli rewards:

$$d_{\text{BER}}(p,q) = p \log \frac{p}{q} + (1-p) \log \frac{1-p}{1-q}$$

• for exponential rewards $p_{\theta_a}(x) = \theta_a e^{-\theta_a x}$:

$$d_{\exp}(u,v) = u - v + u \log \frac{u}{v}$$

The analysis is generic and yields a non-asymptotic regret bound optimal in the sense of Lai and Robbins.

Regret bound

Theorem: Assume that all arms belong to a canonical, regular, exponential family $\mathcal{F} = \{\nu_{\theta} : \theta \in \Theta\}$ of probability distributions indexed by its natural parameter space $\Theta \subseteq \mathbb{R}$. Then, with the choice $f(t) = \log(t) + 3\log\log(t)$ for $t \ge 3$, the number of draws of any suboptimal arm *a* is upper bounded for any horizon $T \ge 3$ as

$$\mathbb{E}[N_{a}(T)] \leq \frac{\log(T)}{d(\mu_{a},\mu^{\star})} + 2\sqrt{\frac{2\pi\sigma_{a,\star}^{2}\left(d'(\mu_{a},\mu^{\star})\right)^{2}}{\left(d(\mu_{a},\mu^{\star})\right)^{3}}}\sqrt{\log(T) + 3\log(\log(T))} + \left(4e + \frac{3}{d(\mu_{a},\mu^{\star})}\right)\log(\log(T)) + 8\sigma_{a,\star}^{2}\left(\frac{d'(\mu_{a},\mu^{\star})}{d(\mu_{a},\mu^{\star})}\right)^{2} + 6,$$

where $\sigma_{a,\star}^2 = \max \{ \operatorname{Var}(\nu_{\theta}) : \mu_a \leq E(\nu_{\theta}) \leq \mu^* \}$ and where $d'(\cdot, \mu^*)$ denotes the derivative of $d(\cdot, \mu^*)$.

Results: Two-Arm Scenario



Figure: Performance of various algorithms when $\theta = (0.9, 0.8)$. Left: average number of draws of the sub-optimal arm as a function of time. Right: box-and-whiskers plot for the number of draws of the sub-optimal arm at time T = 5,000. Results based on 50,000independent replications

Non-parametric setting

Rewards are only assumed to be bounded (say in [0,1])

Need for an estimation procedure

- with non-asymptotic guarantees
- efficient in the sense of Stein / Bahadur
- \implies Idea 1: use d_{BER} (Hoeffding)
- → Idea 2: Empirical Likelihood [Owen '01]
 - Bad idea: use Bernstein / Bennett

First idea: use $d_{\text{\tiny BER}}$

Idea: rescale to [0,1], and take the divergence $d_{\text{\tiny BER}}$.

 \rightarrow because Bernoulli distributions maximize deviations among bounded variables with given expectation:

Lemma (Hoeffding '63)

Let *X* denote a random variable such that $0 \le X \le 1$ and denote by $\mu = \mathbb{E}[X]$ its mean. Then, for any $\lambda \in \mathbb{R}$,

 $E\left[\exp(\lambda X)\right] \le 1 - \mu + \mu \exp(\lambda) .$

This fact is well-known for the variance, but also true for all exponential moments and thus for Cramer-type deviation bounds

Regret Bound for kI-UCB

Theorem

With the divergence d_{BER} , for all T > 3,

$$\mathbb{E}[N_{a}(T)] \leq \frac{\log(T)}{d_{\mathsf{BER}}(\mu_{a},\mu^{\star})} + \frac{\sqrt{2\pi}\log\left(\frac{\mu^{\star}(1-\mu_{a})}{\mu_{a}(1-\mu^{\star})}\right)}{\left(d_{\mathsf{BER}}(\mu_{a},\mu^{\star})\right)^{3/2}} \sqrt{\log(T) + 3\log(\log(T))} + \left(\frac{4e + \frac{3}{d_{\mathsf{BER}}(\mu_{a},\mu^{\star})}}{\left(d_{\mathsf{BER}}(\mu_{a},\mu^{\star})\right)^{2}} + \frac{2\left(\log\left(\frac{\mu^{\star}(1-\mu_{a})}{\mu_{a}(1-\mu^{\star})}\right)\right)^{2}}{\left(d_{\mathsf{BER}}(\mu_{a},\mu^{\star})\right)^{2}} + 6.$$

 kI-UCB satisfies an improved logarithmic finite-time regret bound

Besides, it is asymptotically optimal in the Bernoulli case

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Comparison to UCB

KL-UCB addresses exactly the same problem as UCB, with the same generality, but it has always a smaller regret as can be seen from Pinsker's inequality



Idea 2: Empirical Likelihood

$$U(\hat{\nu}_n, \epsilon) = \sup \left\{ E(\nu') : \nu' \in \mathfrak{M}_1(\operatorname{Supp}(\hat{\nu}_n)) \text{ and } \operatorname{KL}(\hat{\nu}_n, \nu') \le \epsilon \right\}$$

or, rather, modified Empirical Likelihood:

 $U(\hat{\nu}_n, \epsilon) = \sup \left\{ E(\nu') : \nu' \in \mathfrak{M}_1(\operatorname{Supp}(\hat{\nu}_n) \cup \{1\}) \text{ and } \operatorname{KL}(\hat{\nu}_n, \nu') \le \epsilon \right\}$



Coverage properties of the modified EL confidence bound

Proposition: Let $\nu_0 \in \mathfrak{M}_1([0,1])$ with $E(\nu_0) \in (0,1)$ and let X_1, \ldots, X_n be independent random variables with common distribution $\nu_0 \in \mathfrak{M}_1([0,1])$, not necessarily with finite support. Then, for all $\epsilon > 0$,

$$\mathbb{P}\left\{U(\hat{\nu}_n, \epsilon) \le E(\nu_0)\right\} \le \mathbb{P}\left\{K_{inf}(\hat{\nu}_n, E(\nu_0)) \ge \epsilon\right\}$$
$$\le e(n+2)\exp(-n\epsilon) .$$

Remark: For $\{0,1\}$ -valued observations, it is readily seen that $U(\hat{\nu}_n, \epsilon)$ boils down to the upper-confidence bound above. \implies This proposition is at least not always optimal: the presence of the factor n in front of the exponential $\exp(-n\epsilon)$ term is questionable.

Regret bound

Theorem: Assume that \mathcal{F} is the set of finitely supported probability distributions over S = [0, 1], that $\mu_a > 0$ for all arms a and that $\mu^* < 1$. There exists a constant $M(\nu_a, \mu^*) > 0$ only depending on ν_a and μ^* such that, with the choice $f(t) = \log(t) + \log(\log(t))$ for $t \ge 2$, for all $T \ge 3$:

$$\mathbb{E}[N_{a}(T)] \leq \frac{\log(T)}{K_{inf}(\nu_{a},\mu^{\star})} + \frac{36}{(\mu^{\star})^{4}} (\log(T))^{4/5} \log(\log(T)) \\ + \left(\frac{72}{(\mu^{\star})^{4}} + \frac{2\mu^{\star}}{(1-\mu^{\star}) K_{inf}(\nu_{a},\mu^{\star})^{2}}\right) (\log(T))^{4/5} \\ + \frac{(1-\mu^{\star})^{2} M(\nu_{a},\mu^{\star})}{2(\mu^{\star})^{2}} (\log(T))^{2/5} \\ + \frac{\log(\log(T))}{K_{inf}(\nu_{a},\mu^{\star})} + \frac{2\mu^{\star}}{(1-\mu^{\star}) K_{inf}(\nu_{a},\mu^{\star})^{2}} + 4.$$

Classical Bandits

Example: truncated Poisson rewards

for each arm 1 ≤ a ≤ 6 is associated with ν_a, a Poisson distribution with expectation (2 + a)/4, truncated at 10.
 N = 10,000 Monte-Carlo replications on an horizon of T = 20,000 steps.



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Take-home message on classical bandit algorithms

- 1 Use kI-UCB rather than UCB-1 or UCB-2
- 2 Use KL-UCB if speed is not a problem
- 3 todo: improve on the deviation bounds, address general non-parametric families of distributions

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- 4 Alternative: Bayesian-flavored methods:
 - Bayes-UCB [Kaufmann, Cappé, G.]
 - Thompson sampling [Kaufmann & al.]

Roadmap

1 Classical Bandits

2 Best arm identification in two-armed bandits

- Lower bounds on the complexities
- The complexity of A/B Testing with Gaussian feedback

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The complexity of A/B Testing with binary feedback

3 Optimal Exploration with Probabilistic Expert Advice

- Missing mass and Good-UCB
- Analysis: Classical and Macroscopic Optimality

Motivation

On the Complexity of Best Arm Identification in Multi-Armed Bandit Models, ArXiv (COLT 2014) joint work with O. Cappé and E. Kaufmann





Improve performance:

➔ fixed number of test users – > smaller probability of error

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➔ fixed probability of error – > fewer test users

Tools: sequential allocation and stopping

A two-armed bandit model is

- a set $\nu = (\nu_1, \nu_2)$ of two probability distributions ('arms') with respective means μ_1 and μ_2
- $a^* = \operatorname{argmax}_a \mu_a$ is the (unknown) best am

To find the best arm, an agent interacts with the bandit model with

- a sampling rule $(A_t)_{t \in \mathbb{N}}$ where $A_t \in \{1, 2\}$ is the arm chosen at time *t* (based on past observations) – > a sample $Z_t \sim \nu_{A_t}$ is observed
- a stopping rule τ indicating when he stops sampling the arms
- a *recommendation rule* $\hat{a}_{\tau} \in \{1, 2\}$ indicating which arm he thinks is best (at the end of the interaction)

In classical A/B Testing, the sampling rule A_t is uniform on $\{1,2\}$ and the stopping rule $\tau = t$ is fixed in advance.

Two possible goals

The agent's goal is to design a strategy $\mathcal{A} = ((A_t), \tau, \hat{a}_{\tau})$ satisfying

Fixed-budget setting	Fixed-confidence setting
au = t	$\mathbb{P}_{\nu}(\hat{a}_{\tau} \neq a^*) \leq \delta$
$p_t(\nu) \coloneqq \mathbb{P}_{\nu}(\hat{a}_t \neq a^*)$ as small as possible	$\mathbb{E}_{ u}[au]$ as small as possible

An algorithm using uniform sampling is

Fixed-budget setting	Fixed-confidence setting
a classical test of	a sequential test of
$(\mu_1 > \mu_2)$ against $(\mu_1 < \mu_2)$	$(\mu_1 > \mu_2)$ against $(\mu_1 < \mu_2)$
based on t samples	with probability of error
	uniformly bounded by δ

[Siegmund 85]: sequential tests can save samples, !

The complexities of best-arm identification

For a class \mathcal{M} bandit models, algorithm $\mathcal{A} = ((A_t), \tau, \hat{a}_{\tau})$ is...

Fixed-budget setting	Fixed-confidence setting	
consistent on $\mathcal M$ if	$\delta extsf{-PAC}$ on $\mathcal M$ if	
$\forall \nu \in \mathcal{M}, p_t(\nu) = \mathbb{P}_{\nu}(\hat{a}_t \neq a^*) \xrightarrow[t \to \infty]{} 0$	$\forall \nu \in \mathcal{M}, \ \mathbb{P}_{\nu}(\hat{a}_{\tau} \neq a^*) \leq \delta$	
From the literature		
$p_t(u) \simeq \exp\left(-rac{t}{CH(u)} ight)$	$\mathbb{E}_{\nu}[\tau] \simeq C' H'(\nu) \log rac{1}{\delta}$	
[Audibert et al. 10],[Bubeck et al. 11]	[Mannor Tsitsilis 04],[Even-Dar al. 06]	
[Bubeck et al. 13],	[Kalanakrishnan et al.12],	

Two complexities

$$\kappa_{\mathsf{B}}(\nu) = \inf_{\mathcal{A} \text{ cons.}} \left(\limsup_{t \to \infty} -\frac{1}{t} \log p_t(\nu) \right)^{-1} \quad \kappa_{\mathsf{C}}(\nu) = \inf_{\mathcal{A} \text{ } \delta -\mathsf{PAC}} \limsup_{\delta \to 0} \frac{\mathbb{E}_{\nu}[\tau]}{\log(1/\delta)}$$

for a probability of error $\leq \delta$,
budget $t \simeq \kappa_B(\nu) \log \frac{1}{\delta}$ for a probability of error $\leq \delta$
 $\mathbb{E}_{\nu}[\tau] \simeq \kappa_C(\nu) \log \frac{1}{\delta}$

Changes of distribution

New formulation for a change of distribution

Let ν and ν' be two bandit models. Let N_1 (resp. N_2) denote the total number of draws of arm 1 (resp. arm 2) by algorithm \mathcal{A}). For any $A \in \mathcal{F}_{\tau}$ such that $0 < \mathbb{P}_{\nu}(A) < 1$

 $\mathbb{E}_{\nu}[N_1]\mathsf{KL}(\nu_1,\nu_1') + \mathbb{E}_{\nu}[N_2]\mathsf{KL}(\nu_2,\nu_2') \ge d_{\mathsf{ber}}\big(\mathbb{P}_{\nu}(A),\mathbb{P}_{\nu'}(A)\big),$

where
$$d_{\text{BER}}(x,y) = x \log(x/y) + (1-x) \log((1-x)/(1-y)).$$

General lower bounds

Theorem 1

Let \mathcal{M} be a class of two armed bandit models that are continuously parametrized by their means. Let $\nu = (\nu_1, \nu_2) \in \mathcal{M}$.

Fixed-budget setting	Fixed-confidence setting	
any consistent algorithm satisfies	any δ -PAC algorithm satisfies	
$\limsup_{t\to\infty} -\frac{1}{t}\log p_t(\nu) \leq K^*(\nu_1,\nu_2)$	$\mathbb{E}_{\nu}[\tau] \geq \frac{1}{K_{*}(\nu_{1},\nu_{2})} \log\left(\frac{1}{2\delta}\right)$	
with $K^*(u_1, u_2)$	with $K_*(u_1, u_2)$	
= $KL(u^*, u_1)$ = $KL(u^*, u_2)$	$=KL(\nu_1,\nu_*)=KL(\nu_2,\nu_*)$	
Thus, $\kappa_B(\nu) \ge \frac{1}{K^*(\nu_1,\nu_2)}$	Thus, $\kappa_C(\nu) \ge \frac{1}{K_*(\nu_1,\nu_2)}$	

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Fixed-budget setting

For fixed (known) values $\sigma_1,\sigma_2,$ we consider Gaussian bandit models

 $\mathcal{M} = \left\{ \nu = \left(\mathcal{N}\left(\mu_1, \sigma_1^2\right), \mathcal{N}\left(\mu_2, \sigma_2^2\right) \right) : (\mu_1, \mu_2) \in \mathbb{R}^2, \mu_1 \neq \mu_2 \right\}$

Theorem 1:

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$$\kappa_B(\nu) \ge \frac{2(\sigma_1 + \sigma_2)^2}{(\mu_1 - \mu_2)^2}$$

• A strategy allocating $t_1 = \left[\frac{\sigma_1}{\sigma_1 + \sigma_2}t\right]$ samples to arm 1 and $t_2 = t - t_1$ samples to arm 1, and recommending the empirical best satisfies

$$\liminf_{t \to \infty} -\frac{1}{t} \log p_t(\nu) \ge \frac{(\mu_1 - \mu_2)^2}{2(\sigma_1 + \sigma_2)^2}$$

$$\kappa_B(\nu) = \frac{2(\sigma_1 + \sigma_2)^2}{(\mu_1 - \mu_2)^2}$$

Fixed-confidence setting: Algorithm

The α -Elimination algorithm with exploration rate $\beta(t, \delta)$

- → chooses A_t in order to keep a proportion $N_1(t)/t \simeq \alpha$
- → if $\hat{\mu}_a(t)$ is the empirical mean of rewards obtained from *a* up to time *t*, $\sigma_t^2(\alpha) = \sigma_1^2/[\alpha t] + \sigma_2^2/(t [\alpha t])$,

 $\tau = \inf\left\{t \in \mathbb{N} : |\hat{\mu}_1(t) - \hat{\mu}_2(t)| > \sqrt{2\sigma_t^2(\alpha)\beta(t,\delta)}\right\}$



→ recommends the empirical best arm $\hat{a}_{\tau} = \underset{\alpha \neq \alpha}{\operatorname{argmax}} \hat{\mu}_{a}(\tau)$

Fixed-confidence setting: Results

From Theorem 1:

$$\mathbb{E}_{\nu}[\tau] \geq \frac{2(\sigma_1 + \sigma_2)^2}{(\mu_1 - \mu_2)^2} \log\left(\frac{1}{2\delta}\right)$$

• $\frac{\sigma_1}{\sigma_1 + \sigma_2}$ -Elimination with $\beta(t, \delta) = \log \frac{t}{\delta} + 2\log \log(6t)$ is δ -PAC and

$$\forall \epsilon > 0, \quad \mathbb{E}_{\nu}[\tau] \le (1+\epsilon) \frac{2(\sigma_1 + \sigma_2)^2}{(\mu_1 - \mu_2)^2} \log\left(\frac{1}{2\delta}\right) + \underset{\delta \to 0}{o_{\epsilon}} \left(\log\frac{1}{\delta}\right)$$

$$\kappa_C(\nu) = \frac{2(\sigma_1 + \sigma_2)^2}{(\mu_1 - \mu_2)^2}$$

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Gaussian distributions: Conclusions

For any two fixed values of σ_1 and σ_2 ,

$$\kappa_B(\nu) = \kappa_C(\nu) = \frac{2(\sigma_1 + \sigma_2)^2}{(\mu_1 - \mu_2)^2}$$

If the variances are equal, $\sigma_1 = \sigma_2 = \sigma$,

$$\kappa_B(\nu) = \kappa_C(\nu) = \frac{8\sigma^2}{(\mu_1 - \mu_2)^2}$$

- uniform sampling is optimal only when $\sigma_1 = \sigma_2$
- 1/2-Elimination is δ -PAC for a smaller exploration rate $\beta(t, \delta) \simeq \log(\log(t)/\delta)$

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Lower bounds for Bernoulli bandit models

 $\mathcal{M} = \{ \nu = (\mathcal{B}(\mu_1), \mathcal{B}(\mu_2)) : (\mu_1, \mu_2) \in]0; 1[^2, \mu_1 \neq \mu_2 \},\$

shorthand: $K(\mu, \mu') = KL(\mathcal{B}(\mu), \mathcal{B}(\mu')).$

Fixed-budget setting	Fixed-confidence setting	
any consistent algorithm satisfies	any δ -PAC algorithm satisfies	
$\limsup_{t\to\infty} -\frac{1}{t}\log p_t(\nu) \leq K^*(\mu_1,\mu_2)$	$\mathbb{E}_{\nu}[\tau] \ge \frac{1}{K_{*}(\mu_{1},\mu_{2})} \log\left(\frac{1}{2\delta}\right)$	
(Chernoff information)		

 $\mathsf{K}^*(\mu_1,\mu_2) > \mathsf{K}_*(\mu_1,\mu_2)$

Algorithms using uniform sampling

	For any consistent	For any δ -PAC
algorithm	$p_t(\nu) \gtrsim e^{-K^*(\mu_1,\mu_2)t}$	$\frac{\mathbb{E}_{\nu}[\tau]}{\log(1/\delta)} \gtrsim \frac{1}{K_{*}(\mu_{1},\mu_{2})}$
algorithm using uniform sampling	$p_t(\nu) \gtrsim e^{-\frac{K(\overline{\mu},\mu_1) + K(\overline{\mu},\mu_2)}{2}t}$ with $\overline{\mu} = f(\mu_1,\mu_2)$	$\frac{\mathbb{E}_{\nu}[\tau]}{\log(1/\delta)} \gtrsim \frac{2}{K(\mu_{1},\mu) + K(\mu_{2},\mu)}$ with $\underline{\mu} = \frac{\overline{\mu_{1} + \mu_{2}}}{2}$

<u>Remark:</u> Quantities in the same column appear to be close from one another

⇒ Binary rewards: uniform sampling close to optimal

Algorithms using uniform sampling

	For any consistent	For any δ -PAC
algorithm	$p_t(\nu) \simeq e^{-K^*(\mu_1,\mu_2)t}$	$\frac{\mathbb{E}_{\nu}[\tau]}{\log(1/\delta)} \gtrsim \frac{1}{K_{*}(\mu_{1},\mu_{2})}$
algorithm using uniform sampling	$p_t(\nu) \simeq e^{-\frac{K(\overline{\mu},\mu_1) + K(\overline{\mu},\mu_2)}{2}t}$ with $\overline{\mu} = f(\mu_1,\mu_2)$	$\frac{\mathbb{E}_{\nu}[\tau]}{\log(1/\delta)} \gtrsim \frac{2}{K(\mu_{1},\mu) + K(\mu_{2},\mu)}$ with $\underline{\mu} = \frac{\overline{\mu_{1} + \mu_{2}}}{2}$

<u>Remark:</u> Quantities in the same column appear to be close from one another

⇒ Binary rewards: uniform sampling close to optimal

Fixed-budget setting

We show that

$$\kappa_B(\nu) = \frac{1}{\mathsf{K}^*(\mu_1, \mu_2)}$$

(matching algorithm not implementable in practice)

The algorithm using uniform sampling and recommending the empirical best arm is preferable (and very close to optimal)

Fixed-confidence setting

 $\delta\text{-PAC}$ algorithms using uniform sampling satisfy

$$\frac{\mathbb{E}_{\nu}[\tau]}{\log(1/\delta)} \geq \frac{1}{I_{*}(\nu)} \text{ with } I_{*}(\nu) = \frac{\mathsf{K}(\mu_{1}, \frac{\mu_{1}+\mu_{2}}{2}) + \mathsf{K}(\mu_{2}, \frac{\mu_{1}+\mu_{2}}{2})}{2}.$$

The algorithm using uniform sampling and

$$\tau = \inf\left\{t \in 2\mathbb{N}^* : |\hat{\mu}_1(t) - \hat{\mu}_2(t)| > \log\frac{\log(t) + 1}{\delta}\right\}$$

is δ -PAC but not optimal: $\frac{\mathbb{E}[\tau]}{\log(1/\delta)} \simeq \frac{2}{(\mu_1 - \mu_2)^2} > \frac{1}{I_*(\nu)}.$

A better stopping rule NOT based on the difference of empirical means

$$\tau = \inf\left\{t \in 2\mathbb{N}^* : tI_*(\hat{\mu}_1(t), \hat{\mu}_2(t)) > \log\frac{\log(t) + 1}{\delta}\right\}$$

Bernoulli distributions: Conclusion

Regarding the complexities:

•
$$\kappa_B(\nu) = \frac{1}{\mathsf{K}^*(\mu_1, \mu_2)}$$

• $\kappa_C(\nu) \ge \frac{1}{\mathsf{K}_*(\mu_1, \mu_2)} > \frac{1}{\mathsf{K}^*(\mu_1, \mu_2)}$

Thus

$$\kappa_C(\nu) > \kappa_B(\nu)$$

Regarding the algorithms

- There is not much to gain by departing from uniform sampling
- In the fixed-confidence setting, a sequential test based on the difference of the empirical means is no longer optimal

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Conclusion on Best Arm Identification

- the complexities $\kappa_B(\nu)$ and $\kappa_C(\nu)$ are not always equal (and feature some different informational quantities)
- for Bernoulli distributions and Gaussian with similar variances, strategies using uniform sampling are (almost) optimal
- strategies using random stopping do not necessarily lead to a saving in terms of the number of sample used
- Generalization to *m* best arms identification among *K* arms

Roadmap

1 Classical Bandits

- 2 Best arm identification in two-armed bandits
 Lower bounds on the complexities
 The complexity of A/B Testing with Gaussian feedback
 The complexity of A/B Testing with binary feedback
- 3 Optimal Exploration with Probabilistic Expert Advice
 - Missing mass and Good-UCB
 - Analysis: Classical and Macroscopic Optimality

Optimal Discovery with Probabilistic Expert Advice: Finite Time Analysis and Macroscopic Optimality, JMLR 2013 joint work with S. Bubeck and D. Ernst

- Subset A ⊂ X of important items
- $\blacksquare |\mathcal{X}| \gg 1, |A| \ll |\mathcal{X}|$
- Access to X only by probabilistic experts (P_i)_{1≤i≤K}: sequential independent draws



Goal: discover rapidly the elements of $A_{(n)}, A_{(n)}, A_{(n)}$

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- Subset A ⊂ X of important items
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Goal: discover rapidly the elements of A

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- Subset *A* ⊂ *X* of important items
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Goal: discover rapidly the elements of A

Optimal Exploration with Probabilistic Expert Advice

Search space : $A \in \Omega$ discrete set Probabilistic experts : $P_i \in \mathfrak{M}_1(\Omega)$ for $i \in \{1, \dots, K\}$ Requests : at time t, calling expert I_t yields a realization of $X_t = X_{I_t,t}$ independent with law P_a

Goal : find as many distinct elements of *A* as possible with few requests :

$$F_n = \operatorname{Card} \left(A \cap \{ X_1, \dots, X_n \} \right)$$

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Goal

At each time step t = 1, 2, ...:

- pick an index $I_t = \pi_t (I_1, Y_1, \dots, I_{s-1}, Y_{s-1}) \in \{1, \dots, K\}$ according to past observations
- observe $Y_t = X_{I_t, n_{I_t, t}} \sim P_{I_t}$, where

$$n_{i,t} = \sum_{s \le t} \mathbb{1}\{I_s = i\}$$

Goal: design the strategy $\pi = (\pi_t)_t$ so as to maximize the number of important items found after *t* requests

$$F^{\pi}(t) = \left| A \cap \left\{ Y_1, \dots, Y_t \right\} \right|$$

Assumption: non-intersecting supports

 $A \cap \operatorname{supp}(P_i) \cap \operatorname{supp}(P_j) = \emptyset$ for $i \neq j$

Is it a Bandit Problem ?

It looks like a bandit problem...

- sequential choices among K options
- want to maximize cumulative rewards
- exploration vs exploitation dilemma

... but it is not a bandit problem !

- rewards are not i.i.d.
- destructive rewards: no interest to observe twice the same important item
- all strategies eventually equivalent

The oracle strategy

Proposition: Under the non-intersecting support hypothesis, the greedy oracle strategy selecting the expert with highest 'missing mass'

$$I_t^* \in \operatorname*{arg\,max}_{1 \le i \le K} P_i \left(A \smallsetminus \{Y_1, \dots, Y_t\} \right)$$

is optimal: for every possible strategy π , $\mathbb{E}[F^{\pi}(t)] \leq \mathbb{E}[F^{*}(t)]$.

Remark: the proposition if false if the supports may intersect

⇒ estimate the "missing mass of important items"!

Missing mass estimation

Let us first focus on one expert *i*: $P = P_i, X_n = X_{i,n}$

 X_1, \ldots, X_n independent draws of P

$$O_n(x) = \sum_{m=1}^n \mathbb{1}\{X_m = x\}$$



How to 'estimate' the total mass of the unseen important items

$$R_n = \sum_{x \in A} P(x) \mathbb{1}\{O_n(x) = 0\}?$$

The Good-Turing Estimator

Idea: use the **hapaxes** = items seen only once (linguistic)

$$\hat{R}_n = \frac{U_n}{n}$$
, where $U_n = \sum_{x \in A} \mathbb{1}\{O_n(x) = 1\}$

Lemma [Good '53]: For every distribution P,

$$0 \le \mathbb{E}[\hat{R}_n] - \mathbb{E}[R_n] \le \frac{1}{n}$$

Proposition: With probability at least $1 - \delta$ for *every P*,

$$\hat{R}_n - \frac{1}{n} - (1 + \sqrt{2})\sqrt{\frac{\log(4/\delta)}{n}} \le R_n \le \hat{R}_n + (1 + \sqrt{2})\sqrt{\frac{\log(4/\delta)}{n}}$$

See [McAllester and Schapire '00, McAllester and Ortiz '03]:
deviations of *R_n*: McDiarmid's inequality
deviations of *R_n*: negative association

The Good-UCB algorithm [Bubeck, Ernst & G.]

Optimistic algorithm based on Good-Turing's estimator :

$$I_{t+1} = \operatorname*{arg\,max}_{i \in \{1, \dots, K\}} \left\{ \frac{H_i(t)}{N_i(t)} + c \sqrt{\frac{\log(t)}{N_i(t)}} \right\}$$

- $N_i(t)$ = number of draws of P_i up to time t
- $H_i(t)$ = number of elements of A seen exactly once thanks to P_i

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c = tuning parameter

Classical analysis

Theorem: For any $t \ge 1$, under the non-intersecting support assumption, Good-UCB (with constant $C = (1 + \sqrt{2})\sqrt{3}$) satisfies

$$\mathbb{E}\left[F^*(t) - F^{UCB}(t)\right] \le 17\sqrt{Kt\log(t)} + 20\sqrt{Kt} + K + K\log(t/K)$$

Remark: Usual result for bandit problem, but not-so-simple analysis

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A Typical Run of Good-UCB



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The macroscopic limit

- **Restricted framework:** $P_i = \mathcal{U}\{1, \dots, N\}$
- $\blacksquare N \to \infty$
- $|A \cap \operatorname{supp}(P_i)|/N \to q_i \in (0,1), q = \sum_i q_i$



The macroscopic limit

- **Restricted framework:** $P_i = \mathcal{U}\{1, \dots, N\}$
- $\blacksquare N \to \infty$
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The macroscopic limit

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The Oracle behaviour

The limiting discovery process of the Oracle strategy is *deterministic*

Proposition: For every $\lambda \in (0, q_1)$, for every sequence $(\lambda^N)_N$ converging to λ as N goes to infinity, almost surely

$$\lim_{N \to \infty} \frac{T^N_*(\lambda^N)}{N} = \sum_i \left(\log \frac{q_i}{\lambda} \right)_+$$

Oracle vs. uniform sampling

Oracle: The proportion of important items not found after *Nt* draws tends to

$$q - F^*(t) = I(t)\underline{q}_{I(t)} \exp\left(-t/I(t)\right) \le K\underline{q}_K \exp\left(-t/K\right)$$

with $\underline{q}_{K} = \left(\prod_{i=1}^{K} q_{i}\right)^{1/K}$ the geometric mean of the $(q_{i})_{i}$.

- Uniform: The proportion of important items not found after Nt draws tends to $K\bar{q}_K \exp(-t/K)$
- → Asymptotic ratio of efficiency

$$\rho(q) = \frac{\overline{q}_K}{\underline{q}_K} = \frac{\frac{1}{K} \sum_{i=1}^k q_i}{\left(\prod_{i=1}^k q_i\right)^{1/K}} \ge 1$$

larger if the $(q_i)_i$ are unbalanced

Macroscopic optimality

Theorem: Take $C = (1 + \sqrt{2})\sqrt{c+2}$ with c > 3/2 in the Good-UCB algorithm.

For every sequence (λ^N)_N converging to λ as N goes to infinity, almost surely

$$\limsup_{N \to +\infty} \frac{T_{UCB}^N(\lambda^N)}{N} \le \sum_i \left(\log \frac{q_i}{\lambda}\right)_+$$

The proportion of items found after Nt steps F^{GUCB} converges uniformly to F* as N goes to infinity

Simulation



Number of items found by Good-UCB (line), the oracle (bold dashed), and by uniform sampling (light dotted) as a function of time, for sample sizes N = 128, N = 500, N = 1000 et N = 10000, in an environment with 7 experts.