

The Sample Complexity of Level Set Approximation

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Problem, motivations, related work

Problem

Approximating $\{\boldsymbol{x} : f(\boldsymbol{x}) = a\} \subset [0, 1]^d$

- ▶ $f : [0, 1]^d \rightarrow \mathbb{R}$ unknown in some known smoothness class
- ▶ $a \in \mathbb{R}$ a fixed known threshold

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Determining parameters that result in a given outcome (computer experiments, uncertainty quantification, nuclear engineering, coastal flooding, etc.)

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Related work

- ▶ **Gaussian process models:** Chevalier et al. 2014 Technometrics, Azzimonti et al. 2020 Technometrics, Gotovos et al. 2013 IJCAI
- ▶ **Global optimization algorithms:** DOO, Munos 2011 NIPS, HOO, Bubeck et al. 2011 JMLR

Protocol and Objective

Online Protocol

For $n = 1, 2, \dots$

1. pick the next query point \mathbf{x}_n
2. observe the value $f(\mathbf{x}_n)$
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Our Goal

Quantifying the **sample complexity**, i.e., smallest number of evaluations of f needed to

$$\{\mathbf{x} : f(\mathbf{x}) = a\} \subset S_n \subset \{\mathbf{x} : |f(\mathbf{x}) - a| \leq \varepsilon\}$$

for some error level $\varepsilon > 0$

A Hard Problem

Definition

The **packing number** of a non-empty set E is

$$\mathcal{N}(E, \varepsilon) := \sup \left\{ k \in \mathbb{N} : \exists \mathbf{x}_1, \dots, \mathbf{x}_k \in E, \min_{i \neq j} \|\mathbf{x}_i - \mathbf{x}_j\|_\infty > \varepsilon \right\}$$

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Theorem

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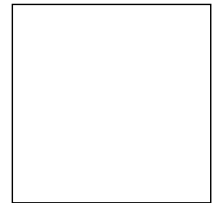
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Not surprising, the level set is defined by a single equation in d unknowns

A General Solution

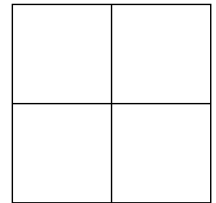
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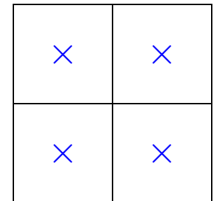
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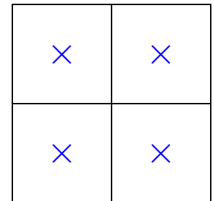
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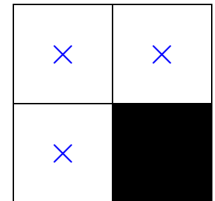
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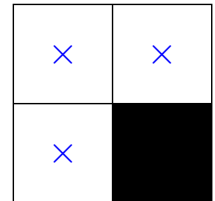
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Theorem

If the g_C 's are "accurate approximations" of f on the C 's,

$$\text{sample complexity of BA} \lesssim \sum_{i=1}^{i(\varepsilon)} \mathcal{N}(\{|f - a| \leq c_i\}, d_i)$$

where $i(\varepsilon) \approx \log(1/\varepsilon)$, $c_1 > c_2 > \dots$, $d_1 > d_2 > \dots$ depend on the g_C 's and their error bounds

Consequence for γ -Hölder functions

γ -Hölder functions

f is γ -Hölder ($|f(\mathbf{x}) - f(\mathbf{y})| \leq c \|\mathbf{x} - \mathbf{y}\|^\gamma$, with $\gamma \in (0, 1]$)

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Local approximator for BAH

- ▶ For a cell C (hypercube), we query the center and take the local approximator g_C as constant
- ▶ The error of g_C on C is $\lesssim \text{Diam}(C)^\gamma$

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Theorem (upper and lower bound)

The worst-case optimal sample complexity is attained by BAH and

$$\text{sample complexity of BAH} \lesssim \frac{1}{\varepsilon^{d/\gamma}}$$

For lower bound counter example functions are “flat + local bump” (classical)

Consequence for functions with γ_1 -Hölder gradients

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∇f is γ_1 -Hölder

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- ▶ For a cell C (hypercube), we query the 2^d vertices and take the local approximator g_C as multilinear interpolating
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Convexity helps a little: $d \mapsto d - 1$

Theorem

When f has γ_1 -Hölder gradient and is **convex** (+ quantitative conditions), then the worst-case optimal sample complexity is attained by BAG and

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Follows from geometric arguments on level sets of convex functions:

Theorem

If f is convex (+ quantitative conditions), there exists a constant $C^* > 0$ such that

$$\forall r \in (0, 1), \mathcal{N}(\{|f - a| \leq r\}, r) \leq C^* \left(\frac{1}{r}\right)^{d-1}$$

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The paper:




F. Bachoc, T. Cesari and S. Gerchinovitz, “The sample complexity of level set approximation” AISTATS 2021 - oral presentation

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*Thank
you!*