

Puzzles

an overview of some results

Pascale Roesch

Institut de Mathématiques de Toulouse

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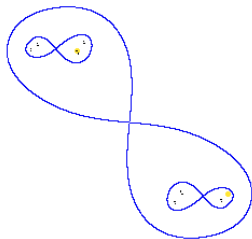
A puzzle is a partition that allows to study long-term dynamics on a space of sequences of abstract symbols

Quadratic Cantor Julia sets

If $\text{degree}(P) = 2$ and if $P^n(0) \rightarrow \infty$, then $J(P)$ is a Cantor set.

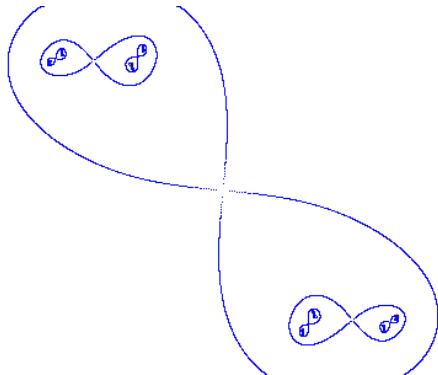
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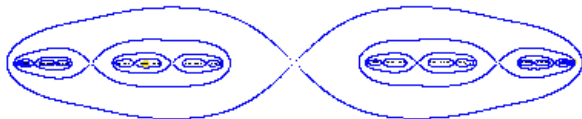
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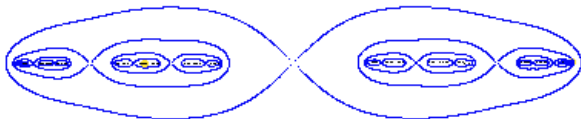
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P is conjugated to the shift σ on two symbols $\Sigma = \{0, 1\}^{\mathbb{N}}$.

Branner-Hubbard Theorem

P. Fatou and G. Julia proved the following theorem.

Theorem A

- 1 The Julia set of a complex polynomial f is connected if and only if $K(f)$ contains all critical points of f .
- 2 The Julia set of a complex polynomial f is a Cantor set if $K(f)$ contains no critical points of f .

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Using combinatorial system of tableaux, Branner and Hubbard completely settled the question of when the Julia set of a cubic polynomial is a Cantor set.

Emergence of the puzzles.

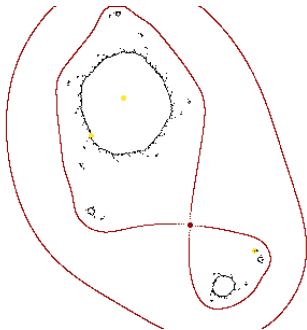
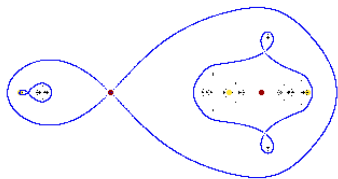
Theorem (B-H)

For a cubic polynomial f with one critical point in $K(f)$, the Julia set $J(f)$ is a Cantor set if and only if the critical components of $K(f)$ are aperiodic.

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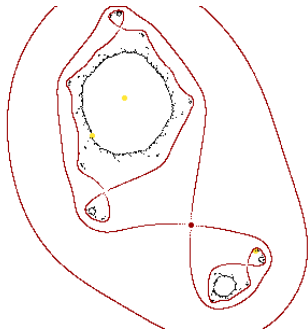
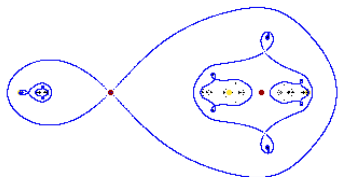
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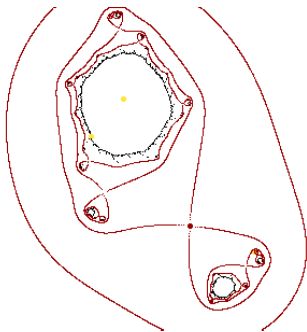
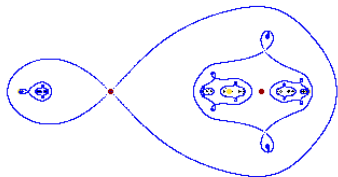
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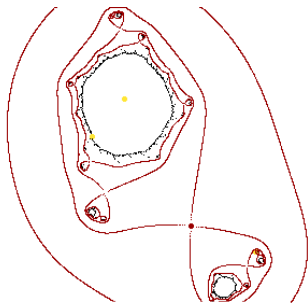
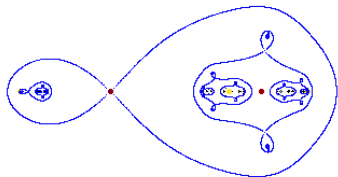
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Theorem (McMullen)

For a cubic polynomial f with Cantor Julia set, the Lebesgue measure of $J(f)$ is zero.

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- $f(P_{n+1}(x)) = P_n(f(x))$
- $f : P_{n+1}(x) \rightarrow P_n(f(x))$ is a covering of degree at most 2

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Every point $x \in K(f)$ defines a "nest" of puzzle pieces

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- $K(f)$ is a Cantor set $\iff \text{diam}(P_n(x)) \rightarrow 0$ for every x .
- K_x is k -periodic \iff the nest is k -periodic :
 $f^k(P_{n+k}(x)) = P_n(x)$ for $n \geq n_0$.

Branner-Hubbard Tableaux

The dynamics can be read on the diagonal of the tableaux

$$P_0(x)$$

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	$f \nearrow$		$f \nearrow$	
$P_2(x)$		$P_2(f(x))$		\vdots
	$f \nearrow$	\vdots	\vdots	\vdots
$P_3(x)$		\vdots	\vdots	
\vdots	\vdots	\vdots	\vdots	$P_{n-2}(f^2(x))$
\vdots	\vdots	$P_{n-1}(f(x))$	$f \nearrow$	$P_{n-1}(f^2(x))$
$P_n(x)$	$f \nearrow$	$P_n(f(x))$	$f \nearrow$	$P_n(f^2(x))$
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\vdots		\vdots		\vdots

Some Analysis

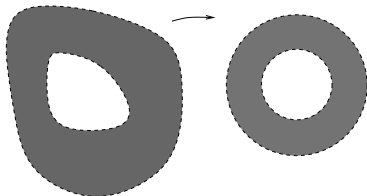
To prove that $K_x = \bigcap P_n(x)$ is reduced to $\{x\}$ one needs to understand this combinatorics and the following analysis.

- 1 The **modulus** of an annulus A estimates its "size", it is a conformal invariant and $\text{mod} (D_R \setminus \overline{D_1}) = \frac{1}{2\pi} \log(R)$;

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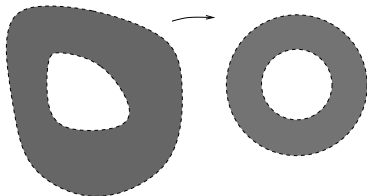


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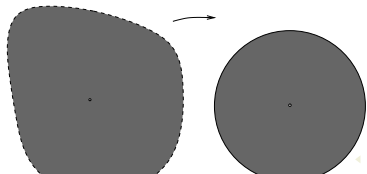
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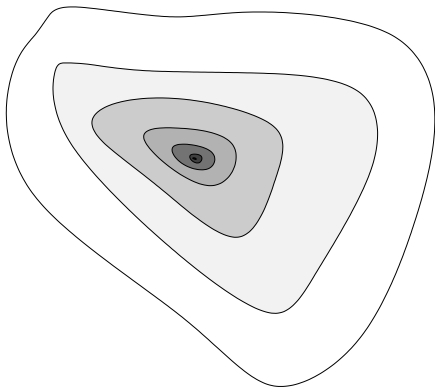
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- 3 it is enough to prove that $\sum \text{mod}(A_n(x)) = \infty$.



Generally $f(A_{n+1}(x)) \neq A_n(f(x))$ for $A_n(x) = P_n(x) \setminus \overline{P_{n+1}(x)}$ but

- It is **critical** if $P_{n+1}(x)$ contains the critical point and $\text{mod}(A_n(x)) = \frac{1}{2} \text{mod}(A_{n-1}(f(x)))$.

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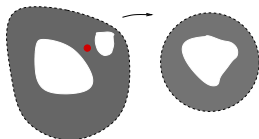
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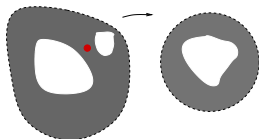


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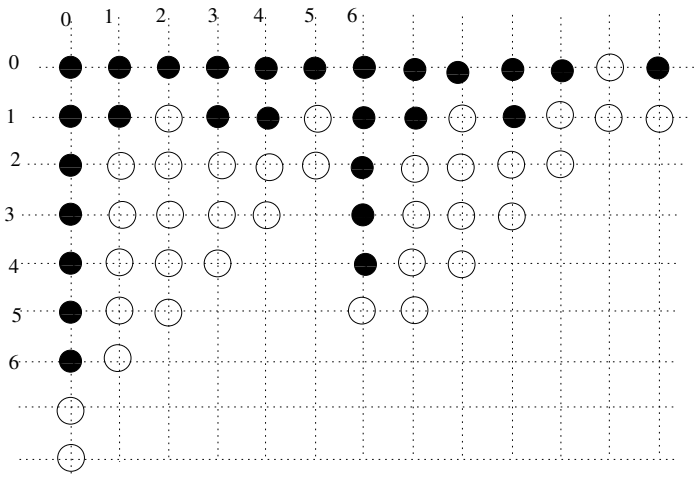


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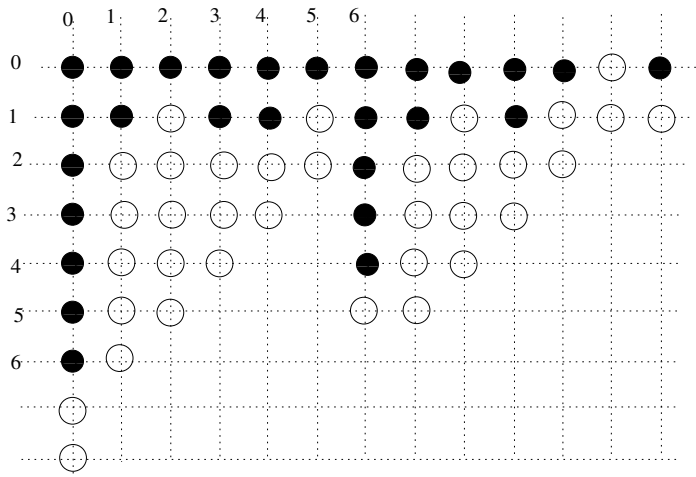


- It is **non critical** if $P_n(x)$ contains no critical point and $\text{mod}(A_n(x)) = \text{mod}(A_{n-1}(f(x)))$.

Tableaux



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Rule 1: $T(i_0, j) = \bullet \implies T(i, j) = \bullet$ for $i \leq i_0$

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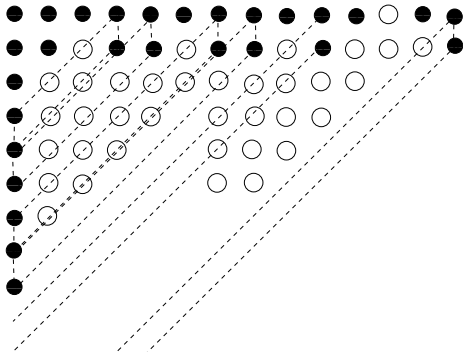
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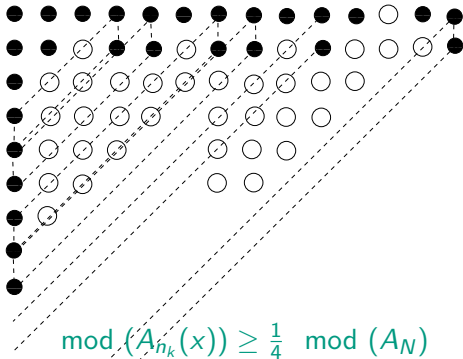
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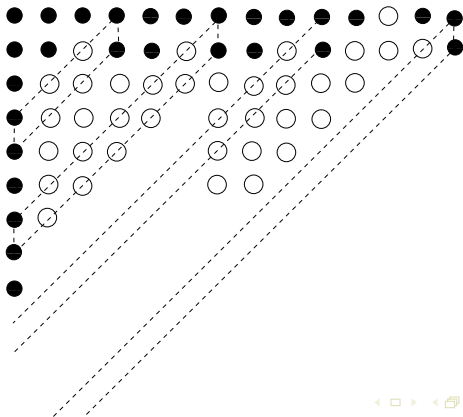
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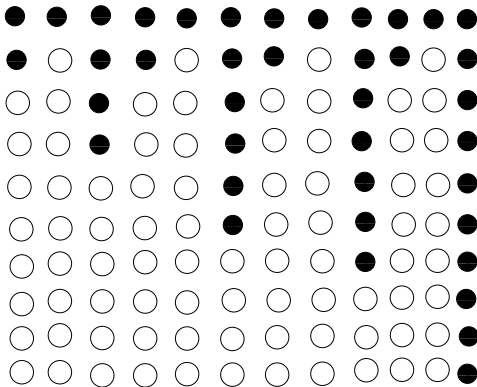
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then there are critical position at any level in the tableau of x

$$\sum \text{ mod } (A_n(x)) \geq \frac{1}{2} \sum \text{ mod } (A_n(c))$$



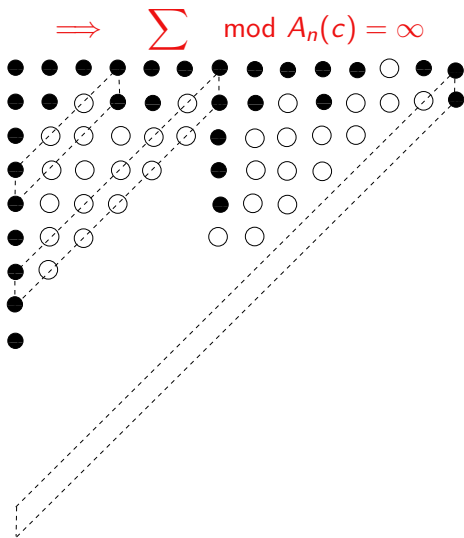
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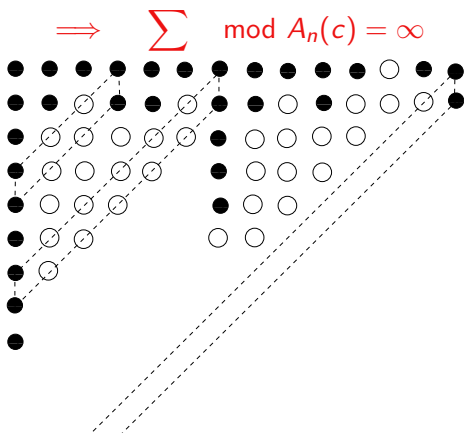
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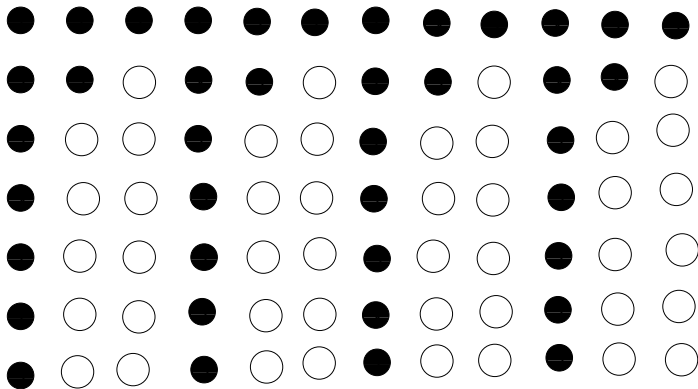
Definition :

An annulus A' is a **child** of an annulus A if there exists some $k > 0$ such that $f^k : A' \rightarrow A$ is a **degree 2 non ramified covering**.

It may happen that no annulus has infinitely many children.

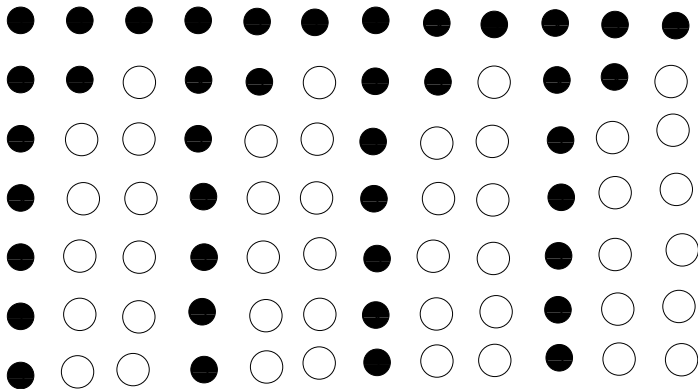
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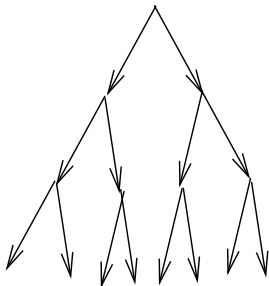
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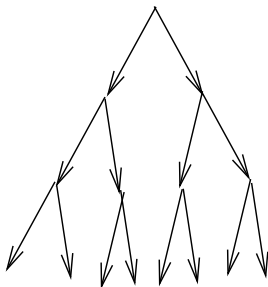


Assume now that the nest ($P_n(c)$) is not periodic

Two is enough

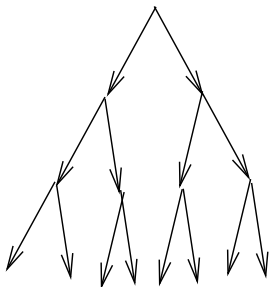


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At least always two children implies $\sum \text{mod } A_n = \infty$

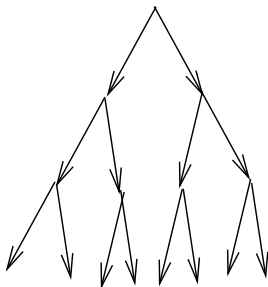
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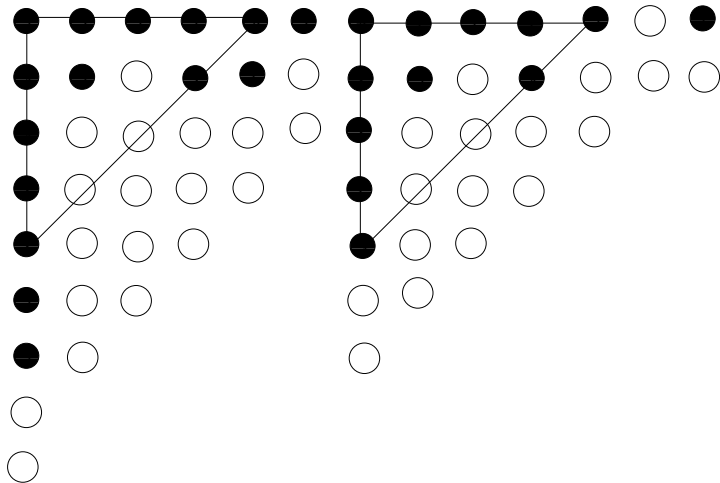
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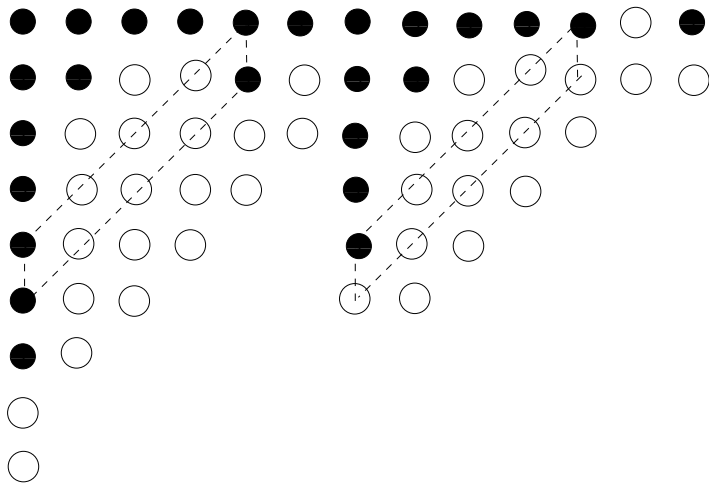
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- 2 **If there is no unique child, every annulus has at least two children and we obtain infinite modulus**

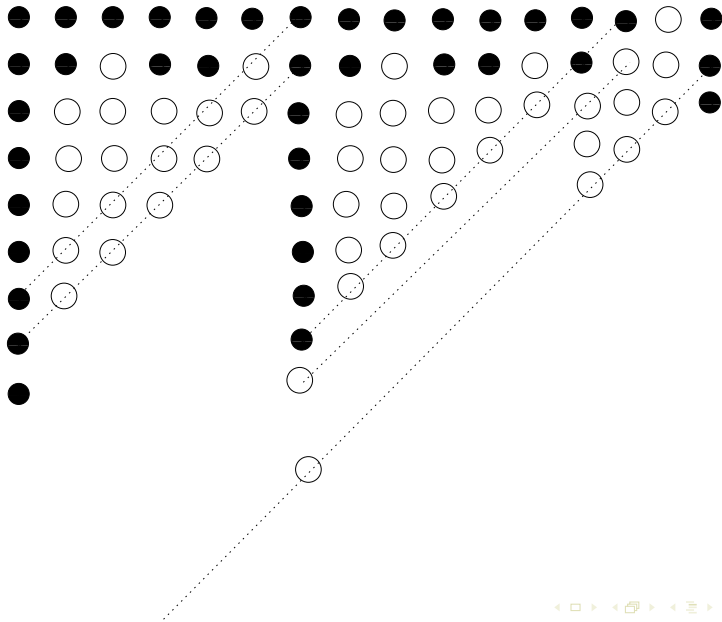
Rule 2



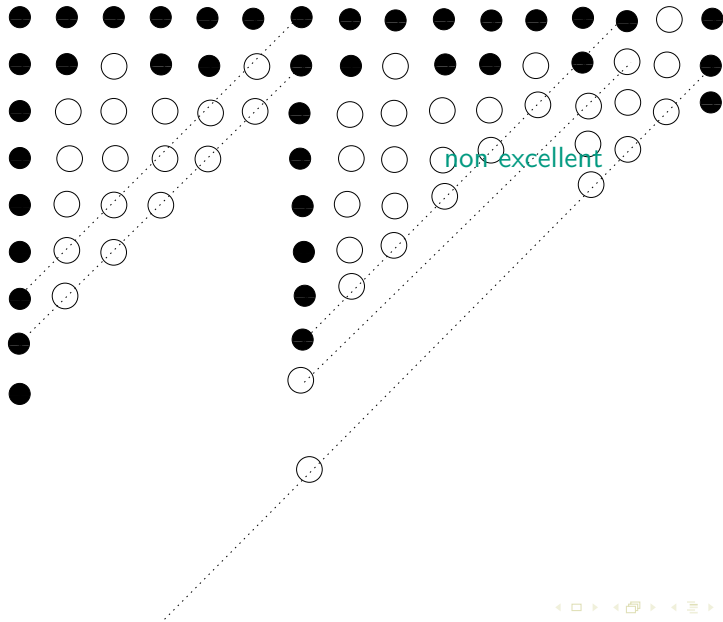
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Proposition :

If the nest $(P_n(c))$ is periodic then K_c is (quasi-conformally) homeomorphic to the filled-in Julia set of a quadratic polynomial.

(by Douady-Hubbard straightening Theorem)

Generalization?

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OF PUZZLES

Key Remark :

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WHAT FOR???

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- 5 get the conjugacy since there are no other identifications than the ones given by the nests shrinking to points.

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Yoccoz puzzle

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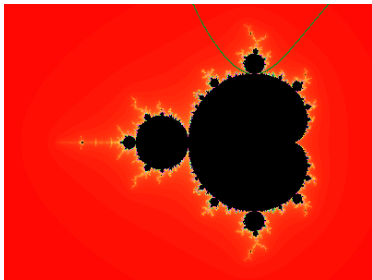
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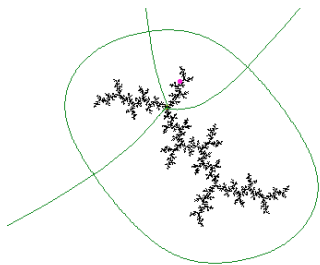
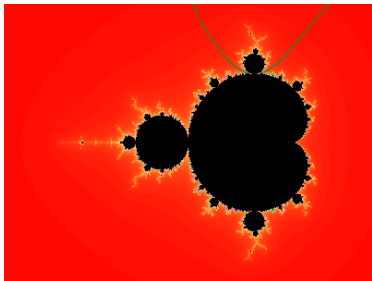
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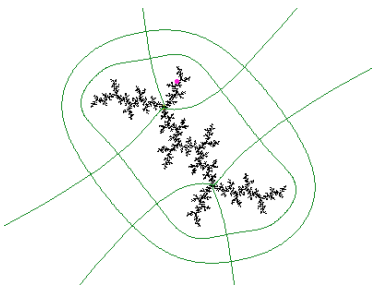
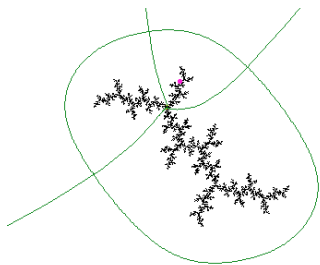
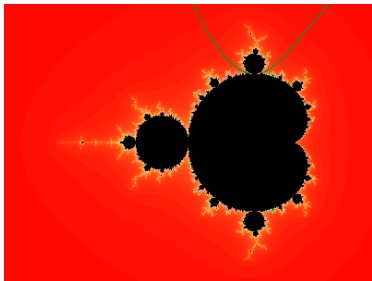
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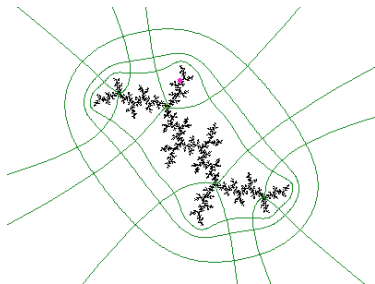
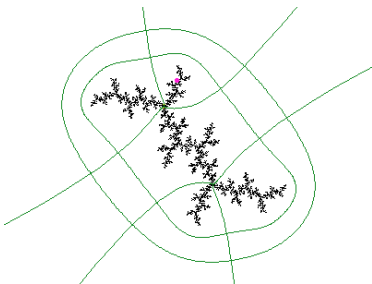
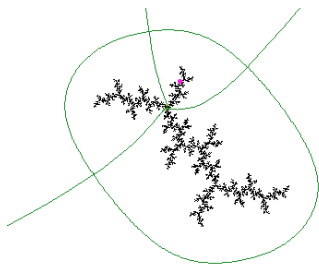
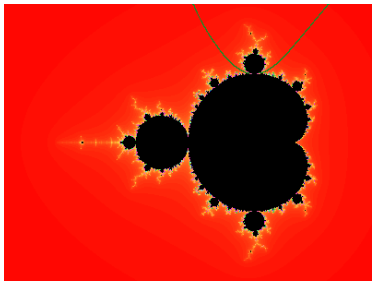
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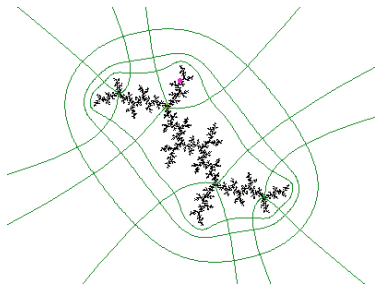
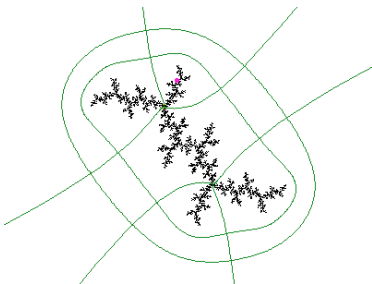
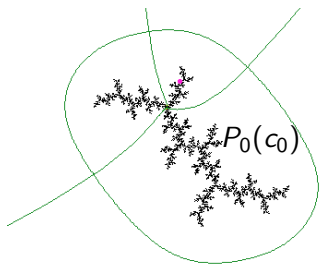
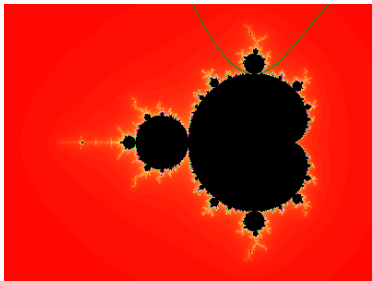
One need to find a good graph

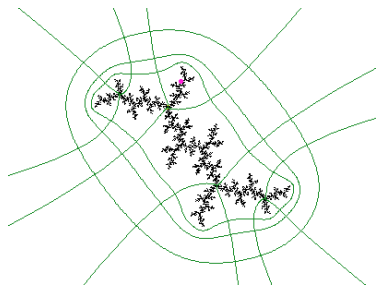
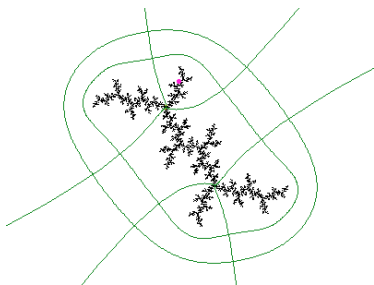
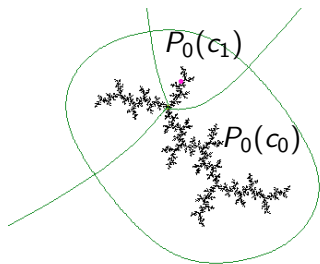
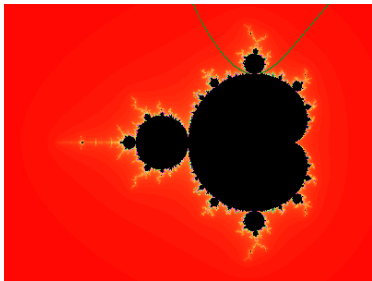


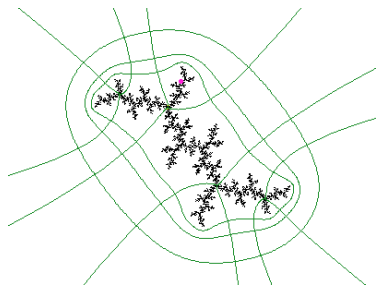
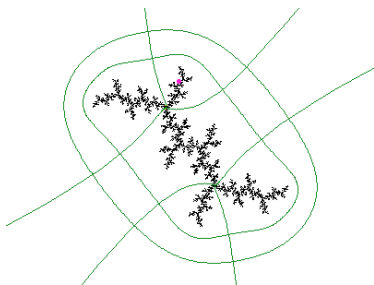
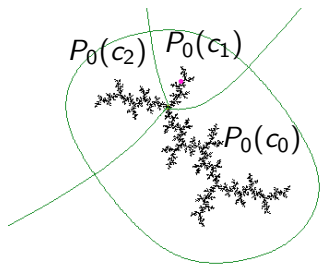
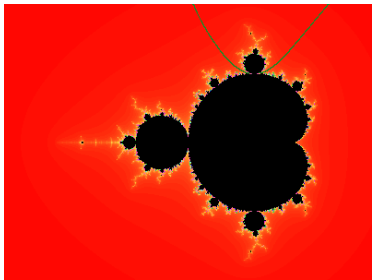


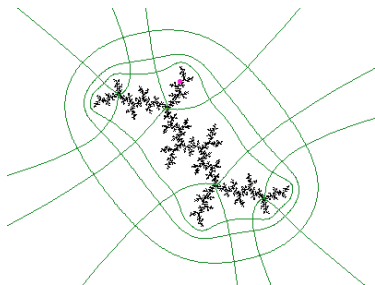
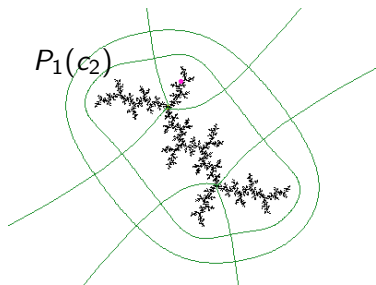
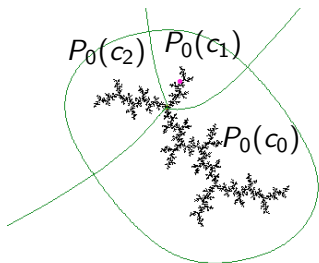
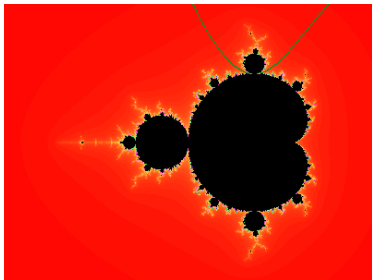


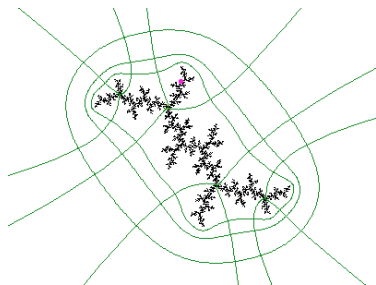
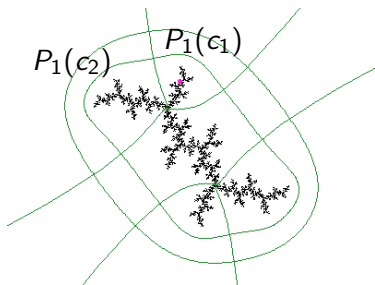
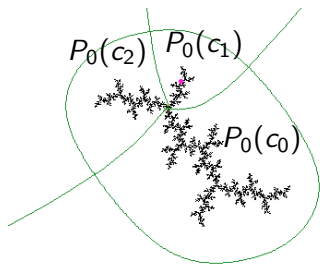
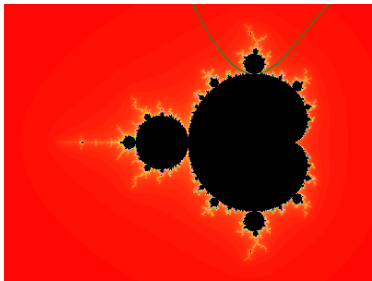


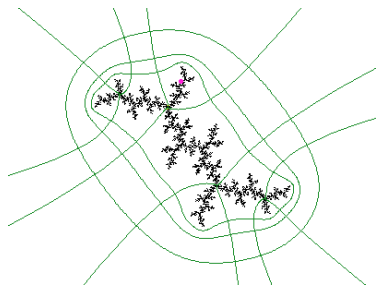
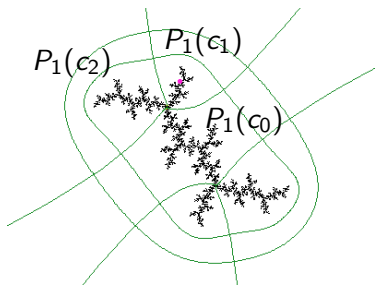
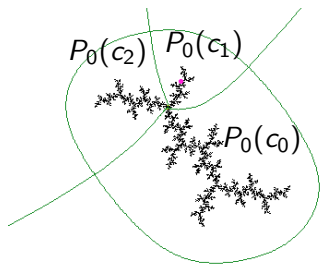
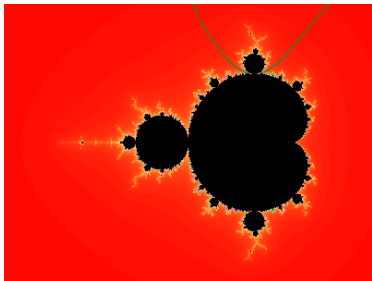


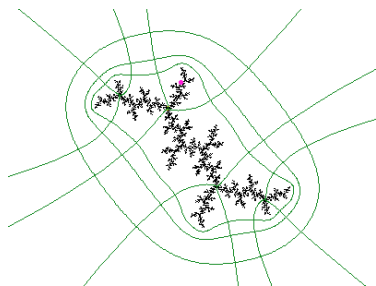
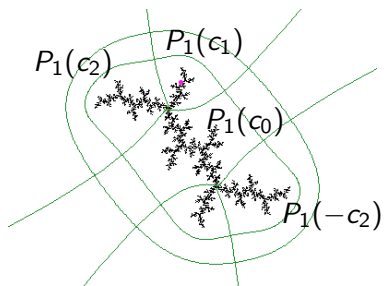
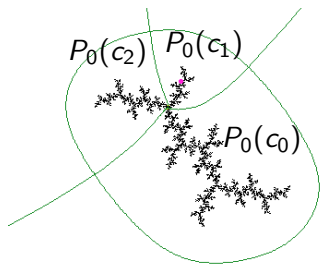
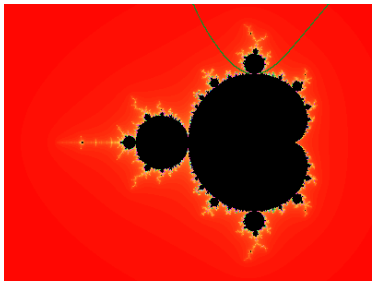


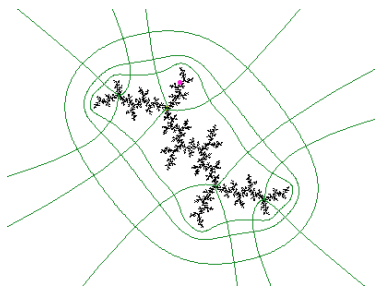
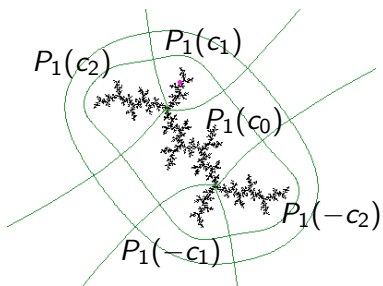
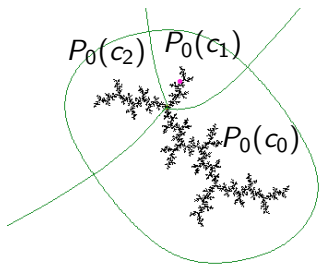
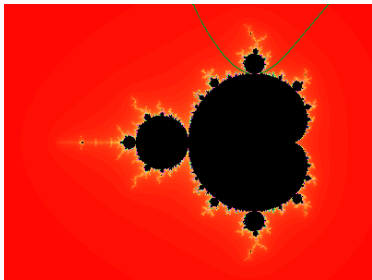








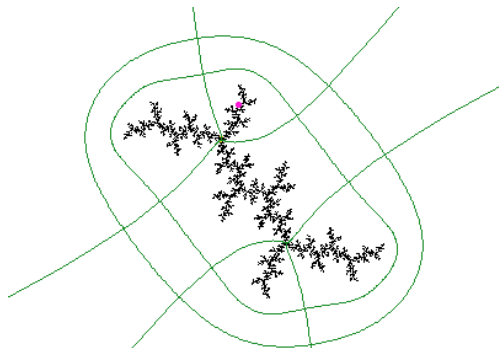




Problem of annuli

A new problem arise here:

TO FIND ENOUGH NON DEGENERATE ANNULI



$P_0(0) \setminus \overline{P_1(-c_1)}$ and $P_0(0) \setminus \overline{P_1(-c_2)}$ are non degenerate.

As before we focus on the critical nest (to be defined we need to)
Assume that the critical orbit does not contain the fixed point α

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$$P_n(f^k(x)) \setminus \overline{P_{n+1}(f^k(x))} \text{ is non degenerate}$$

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Corollary

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The Mandelbrot set is locally connected at such polynomials.

Generalization?

Definition

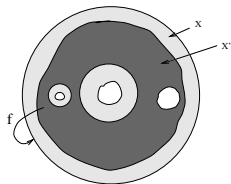
A map $f : X' \rightarrow X$ is (weakly) *rational like* if

- X and X' are connected open sets of $\widehat{\mathbf{C}}$ with smooth boundary, X contains the closure $\overline{X'}$ of X' ($X \supset X'$ and $\partial X \cap \partial X'$ is finite) and ∂X has finitely many connected components;
- $f : X' \rightarrow X$ is a proper holomorphic map with a finite number of critical point and extends continuously to $\overline{X'} \rightarrow \overline{X}$.

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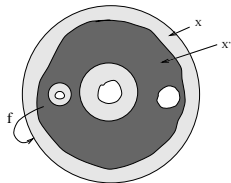
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If X, X' are discs it is polynomial-like.

Definition

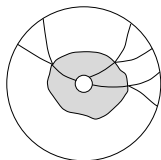
Let $f : X' \rightarrow X$ a (weakly) rational like map with only one critical point, which is simple. A graph Γ is said *admissible* if it satisfies the following conditions:

- Γ is a connected finite graph included in \overline{X} and containing ∂X ;
- Γ is *stable* meaning $f^{-1}(\Gamma)$ contains $\Gamma \cap X'$;
- the forward orbit of the critical point is disjoint from Γ .

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Admissible graph in a basin of attraction

Given an admissible graph Γ for $f : X' \rightarrow X$, the *puzzle pieces of level n* , are the connected components of

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Theorem

Let Γ be an *admissible* graph the that there exists a *non degenerate* annulus surrounding the critical point c .

Let x be a point of $K(f) = \cap f^{-n}(X)$ that is surrounded by infinitely many annuli of Γ , then

- if the critical nest is not periodic, then the puzzles pieces shrink to the point (c or x);
- if the critical nest is periodic, then the map is renormalizable.

Generalization?

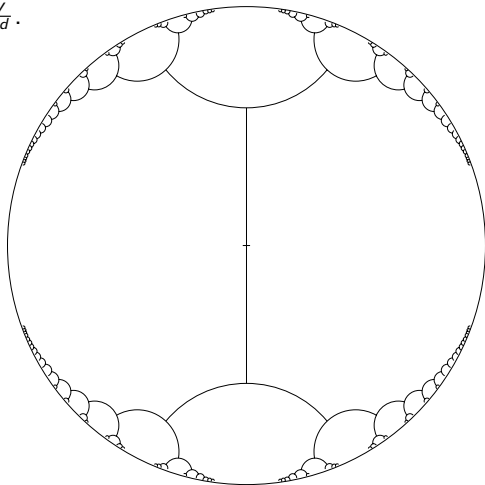
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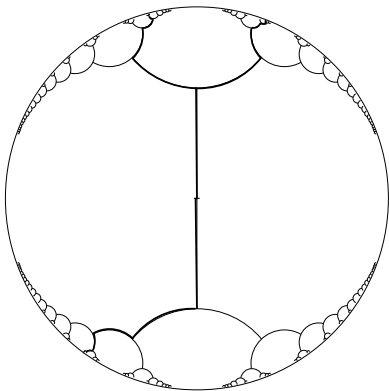
to others kind of graphs

Parabolic rays

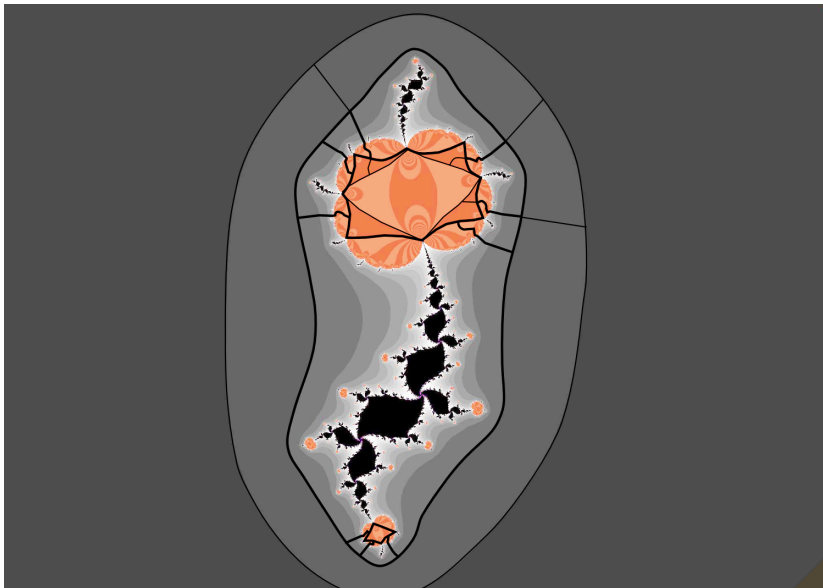
Path of good itinerary in the tree $\mathcal{T} = \cup_n B^{-n}([0, v])$, where

$$B : z \mapsto \frac{z^d + v}{1 + vz^d}.$$





Puzzles



Rational maps

