

1 On the proximity operator of the sum of two closed and convex functions *

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3 **Abstract.** The main result of this paper provides an explicit decomposition of the proximity
4 operator of the sum of two closed and convex functions. For this purpose, we introduce a new oper-
5 ator, called *f-proximity operator*, generalizing the classical notion. After providing some properties
6 and characterizations, we discuss the relations between the *f-proximity operator* and the classical
7 Douglas-Rachford operator. In particular we provide a one-loop algorithm allowing to compute nu-
8 merically this new operator, and thus the proximity operator of the sum of two closed and convex
9 functions. Finally we illustrate the usefulness of our main result in the context of sensitivity analysis
10 of linear variational inequalities of second kind in a Hilbert space.

11 **Key words.** convex analysis, proximity operator, Douglas-Rachford operator, Forward-Back-
12 ward operator, sensitivity analysis, variational inequality.

13 **AMS subject classifications.** 46N10, 47N10, 49J40, 49Q12

14 1. Introduction, notations and basics.

15 **1.1. Introduction.** The *proximity operator* of a proper, closed, convex and
16 extended-real-valued function was first introduced by J.-J. Moreau in 1965 in [18]
17 and can be viewed as an extension of the projection operator on a closed and con-
18 vex subset of a Hilbert space. This wonderful tool plays an important role, from
19 both theoretical and numerical points of view, in convex optimization problems (see,
20 e.g., [5, 16, 20, 22]), inverse problems (see, e.g., [4, 6]), signal processing (see, e.g., [8]),
21 etc. We also refer to [7, 12] and references therein. For the rest of this introduction,
22 we use standard notations of convex analysis. For the reader who is not acquainted
23 with convex analysis, we refer to Section 1.2 for notations and basics.

24 **Motivations from a sensitivity analysis.** The present paper was initially
25 motivated by the sensitivity analysis, with respect to a nonnegative parameter $t \geq 0$, of
26 some parameterized linear variational inequalities of second kind in a Hilbert space H ,
27 with a corresponding functional denoted by $h \in \Gamma_0(H)$. In that framework, the
28 solution $u(t) \in H$ (that depends on the parameter t) can be expressed in terms of
29 the proximity operator of h denoted by prox_h . As a consequence, the differentiability
30 of $u(\cdot)$ at $t = 0$ is strongly related to the regularity of prox_h . If h is a smooth
31 functional, one can easily compute (from the classical implicit function theorem for
32 instance) the differential of prox_h , and then the sensitivity analysis can be achieved.
33 In that smooth case, note that the variational inequality can actually be reduced to an
34 equality. On the other hand, if $h = \iota_K$ is the indicator function of a nonempty closed
35 and convex subset K of H , then $\text{prox}_h = \text{proj}_K$ is the classical projection operator
36 on K . In that case, a result of F. Mignot (see [17, Theorem 2.1 p.145], see also [13,
37 Theorem 2 p.620]) provides an asymptotic development of $\text{prox}_h = \text{proj}_K$ and permits
38 to obtain a differentiability result on $u(\cdot)$ at $t = 0$.

39 In a parallel work (in progress) of the authors on some shape optimization problems

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40 with unilateral contact and friction, the considered variational inequality involves the
 41 sum of two functions, that is, $h = f + g$ where $f = \iota_K$ is the indicator function
 42 of a nonempty closed and convex set of constraints K , and $g \in \Gamma_0(H)$ is a smooth
 43 functional (derived from the regularization of the friction functional in view of a
 44 numerical treatment). Despite the regularity of g , note that the variational inequality
 45 here cannot be reduced to an equality due to the presence of the constraint set K .
 46 In that framework, in order to get an asymptotic development of $\text{prox}_h = \text{prox}_{f+g}$,
 47 a first and natural strategy would be to develop a splitting method, looking for an
 48 explicit expression of prox_{f+g} depending only on the knowledge of prox_f and prox_g .
 49 Unfortunately, this question still remains an open challenge in the literature. Let us
 50 mention that Y.-L. Yu provides in [24] some necessary and/or sufficient conditions on
 51 general functions $f, g \in \Gamma_0(H)$ under which $\text{prox}_{f+g} = \text{prox}_f \circ \text{prox}_g$. Unfortunately,
 52 these conditions are very restrictive and are not satisfied in most of cases.

53 Before coming to the main topic of this paper, we recall that a wide literature is
 54 already concerned with the sensitivity analysis of parameterized (linear and nonlinear)
 55 variational inequalities. We refer for instance to [3, 13, 19, 23] and references therein.
 56 The results in there are considered in very general frameworks. We precise that our
 57 original objective was to look for a simple and compact formula for the derivative $u'(0)$
 58 in the very particular case described above, that is, in the context of a linear variational
 59 inequality and with $h = f + g$ where f is an indicator function and g is a smooth
 60 functional. For this purpose, we were led to consider the proximity operator of the
 61 sum of two proper closed and convex functions, to introduce a new operator and
 62 finally to prove the results presented in this paper.

Introduction of the f -proximity operator and main result. Let us con-
 sider general functions $f, g \in \Gamma_0(H)$. Section 2 is devoted to the introduction (see
 Definition 2.1) of a new operator denoted by prox_g^f , called f -proximity operator of g
 and defined by

$$\text{prox}_g^f := (\text{I} + \partial g \circ \text{prox}_f)^{-1}.$$

63 We prove that its domain satisfies $D(\text{prox}_g^f) = H$ if and only if $\partial(f + g) = \partial f + \partial g$
 64 (see Proposition 2.4), and that prox_g^f can be seen as a generalization of prox_g in
 65 the sense that, if f is constant for instance, then $\text{prox}_g^f = \text{prox}_g$. More general
 66 sufficient (and necessary) conditions under which $\text{prox}_g^f = \text{prox}_g$ are provided in
 67 Propositions 2.11 and 2.14. Note that $\text{prox}_g^f : H \rightrightarrows H$ is *a priori* a set-valued operator.
 68 We provide in Proposition 2.17 some sufficient conditions under which prox_g^f is single-
 69 valued. Some examples illustrate all the previous results throughout the section (see
 70 Examples 2.2, 2.3, 2.6 and 2.16).

71 Finally, if the additivity condition $\partial(f + g) = \partial f + \partial g$ is satisfied, the main result of
 72 the present paper (see Theorem 2.7) provides the equality

$$\text{prox}_{f+g} = \text{prox}_f \circ \text{prox}_g^f.$$

74 Theorem 2.7 allows to prove in a simple and concise way almost all other results of
 75 this paper, making it central in our work.

Relations with the classical Douglas-Rachford operator and algorithms.
 Recall that the proximity operator prox_{f+g} is strongly related to the minimization
 problem

$$\text{argmin } f + g,$$

76 since the solutions are exactly the fixed points of prox_{f+g} . In the sequel, we will
 77 assume that the above problem admits at least one solution. In most of cases, prox_{f+g}
 78 cannot be easily computable, even if prox_f and prox_g are known. As a consequence,
 79 to the best of our knowledge, no *proximal algorithm* $x_{n+1} = \text{prox}_{f+g}(x_n)$, using only
 80 the knowledge of prox_f and prox_g , has been provided in the literature.

The classical *Douglas-Rachford operator*, introduced in [9] and denoted by $\mathcal{T}_{f,g}$ (see
 Section 3 for details), provides an algorithm $x_{n+1} = \mathcal{T}_{f,g}(x_n)$ that is weakly convergent
 to some $x^* \in \text{H}$ satisfying

$$\text{prox}_f(x^*) \in \text{argmin } f + g.$$

81 Even if the *Douglas-Rachford algorithm* is not a proximal algorithm in general, it
 82 is a very powerful tool since it is a one-loop algorithm, allowing to solve the above
 83 minimization problem, that only requires the knowledge of prox_f and prox_g . We refer
 84 to [10, 15] and [2, Section 27.2 p.400] for more details.

85 Section 3 deals with the relations between the Douglas-Rachford operator $\mathcal{T}_{f,g}$ and
 86 the f -proximity operator prox_g^f introduced in this paper. Precisely, for all $x \in \text{H}$,
 87 we prove in Proposition 3.2 that $\text{prox}_g^f(x)$ coincides with the set of fixed points
 88 of $\overline{\mathcal{T}}_{f,g}(x, \cdot)$, where $\overline{\mathcal{T}}_{f,g}(x, \cdot)$ denotes a x -dependent generalization of the classi-
 89 cal Douglas-Rachford operator $\mathcal{T}_{f,g}$. We refer to Section 3 for the precise defini-
 90 tion of $\overline{\mathcal{T}}_{f,g}(x, \cdot)$ that only depends on the knowledge of prox_f and prox_g . In
 91 particular, if $x \in \text{D}(\text{prox}_g^f)$, we prove in Theorem 3.3 that the fixed-point algo-
 92 rithm $y_{k+1} = \overline{\mathcal{T}}_{f,g}(x, y_k)$, denoted by (\mathcal{A}_1) , weakly converges to some $y^* \in \text{prox}_g^f(x)$.

93 Moreover, if the additivity condition $\partial(f + g) = \partial f + \partial g$ is satisfied, we get from
 94 Theorem 2.7 that $\text{prox}_f(y^*) = \text{prox}_{f+g}(x)$. In that situation, we conclude that Algo-
 95 rithm (\mathcal{A}_1) is a one-loop algorithm, that depends only on the knowledge of prox_f and
 96 prox_g , allowing to compute numerically $\text{prox}_{f+g}(x)$.

97 As a consequence, a *proximal-like algorithm* $x_{n+1} = \text{prox}_{f+g}(x_n)$, denoted by (\mathcal{A}_2) ,
 98 using only the knowledge of prox_f and prox_g , can be derived in the above framework
 99 (see Remark 3.7). We refer to Definition 3.5 for the precise meaning of *proximal-like*
 100 *algorithm*.

101 The aim of the present theoretical paper is not to discuss numerical experiments and
 102 comparisons between numerical algorithms (this should be the topic of future works).
 103 However, it should be noted that, in contrary to the classical Douglas-Rachford algo-
 104 rithm, a proximal-like algorithm is a two-loops algorithm. As a consequence, it should
 105 not be expected from Algorithm (\mathcal{A}_2) better performances than the Douglas-Rachford
 106 algorithm for solving the minimization problem $\text{argmin } f + g$.

107 ***Some other applications and forthcoming works.*** Section 4 can be seen as
 108 a conclusion of the paper. Its aim is to provide a glimpse of some other applications
 109 of our main result (Theorem 2.7) and to raise open questions for forthcoming works.
 110 This section is splitted into two parts.

111 In Section 4.1 we consider the framework where $f, g \in \Gamma_0(\text{H})$ with g differentiable
 112 on H . In that framework, we prove from Theorem 2.7 that prox_{f+g} is related to the
 113 classical *Forward-Backward operator* (see [7, Section 10.3 p.191] for details) denoted
 114 by $\mathcal{F}_{f,g}$. Precisely, for all $x \in \text{H}$, we prove in Proposition 4.1 that $\text{prox}_{f+g}(x)$ coincides
 115 with the set of fixed points of $\overline{\mathcal{F}}_{f,g}(x, \cdot)$, where $\overline{\mathcal{F}}_{f,g}(x, \cdot)$ denotes a x -dependent gen-
 116 eralization of the classical Forward-Backward operator $\mathcal{F}_{f,g}$. We refer to Section 4.1

117 for the precise definition of $\overline{\mathcal{F}}_{f,g}(x, \cdot)$ that only depends on the knowledge of prox_f
 118 and ∇g . From this point, one can develop a similar strategy as in Section 3. Precisely,
 119 for all $x \in \mathbb{H}$, one can consider the one-loop algorithm $y_{k+1} = \overline{\mathcal{F}}_{f,g}(x, y_k)$, denoted
 120 by (\mathcal{A}_3) , in order to compute numerically $\text{prox}_{f+g}(x)$, with the only knowledge of
 121 prox_f and ∇g . Moreover, one can also deduce a two-loops algorithm denoted by (\mathcal{A}_4)
 122 as a potential proximal-like algorithm $x_{n+1} = \text{prox}_{f+g}(x_n)$, using only the knowledge
 123 of prox_f and ∇g . Convergence proofs (under some assumptions on f and g) and
 124 numerical experiments of Algorithms (\mathcal{A}_3) and (\mathcal{A}_4) should be the topic of future
 125 works.

In Section 4.2 we return back to our initial motivation, namely the sensitivity analysis, with respect to a nonnegative parameter $t \geq 0$, of some parameterized linear variational inequalities of second kind. Precisely, under some assumptions (see Proposition 4.3 for details), we derive from Theorem 2.7 that if

$$u(t) = \text{prox}_{f+g}(r(t)),$$

where $f = \iota_K$ (where K is a nonempty closed convex set) and where $g \in \Gamma_0(\mathbb{H})$ is a smooth enough functional, then

$$u'(0) = \text{prox}_{\varphi_f + \varphi_g}(r'(0)),$$

126 where $\varphi_f := \iota_C$ (where C is a nonempty closed convex subset of \mathbb{H} related to K)
 127 and where $\varphi_g(x) := \frac{1}{2} \langle D^2 g(u(0))(x), x \rangle$ for all $x \in \mathbb{H}$. We refer to Proposition 4.3
 128 for details. It should be noted that the assumptions of Proposition 4.3 are quite
 129 restrictive, raising open questions about their relaxations (see Remark 4.5). This also
 130 should be the subject of a forthcoming work.

131 **1.2. Notations and basics.** In this section we introduce some notations avail-
 132 able throughout the paper and we recall some basics of convex analysis. We refer to
 133 standard books like [2, 14, 21] and references therein.

Let \mathbb{H} be a real Hilbert space and let $\langle \cdot, \cdot \rangle$ (resp. $\| \cdot \|$) be the corresponding scalar product (resp. norm). For all subset S of \mathbb{H} , we denote respectively by $\text{int}(S)$ and $\text{cl}(S)$ its interior and its closure. In the sequel we denote by $I : \mathbb{H} \rightarrow \mathbb{H}$ the identity application and by $L_x : \mathbb{H} \rightarrow \mathbb{H}$ the affine operator defined by

$$L_x(y) := x - y,$$

134 for all $x, y \in \mathbb{H}$.

For a set-valued map $A : \mathbb{H} \rightrightarrows \mathbb{H}$, the *domain* of A is given by

$$D(A) := \{x \in \mathbb{H} \mid A(x) \neq \emptyset\}.$$

We denote by $A^{-1} : \mathbb{H} \rightrightarrows \mathbb{H}$ the set-valued map defined by

$$A^{-1}(y) := \{x \in \mathbb{H} \mid y \in A(x)\},$$

for all $y \in \mathbb{H}$. Note that $y \in A(x)$ if and only if $x \in A^{-1}(y)$, for all $x, y \in \mathbb{H}$. The *range* of A is given by

$$R(A) := \{y \in \mathbb{H} \mid A^{-1}(y) \neq \emptyset\} = D(A^{-1}).$$

We denote by $\text{Fix}(A)$ the set of all fixed points of A , that is, the set given by

$$\text{Fix}(A) := \{x \in \mathbb{H} \mid x \in A(x)\}.$$

135 Finally, if $A(x)$ is a singleton for all $x \in D(A)$, we say that A is *single-valued*.

For all extended-real-valued functions $g : \mathbb{H} \rightarrow \mathbb{R} \cup \{+\infty\}$, the *domain* of g is given by

$$\text{dom}(g) := \{x \in \mathbb{H} \mid g(x) < +\infty\}.$$

136 We say that g is *proper* if $\text{dom}(g) \neq \emptyset$, and that g is *closed* (or *lower semi-continuous*)
 137 if its epigraph is a closed subset of $\mathbb{H} \times \mathbb{R}$.

Let $g : \mathbb{H} \rightarrow \mathbb{R} \cup \{+\infty\}$ be a proper extended-real-valued function. We denote by $g^* : \mathbb{H} \rightarrow \mathbb{R} \cup \{+\infty\}$ the *conjugate* of g defined by

$$g^*(y) := \sup_{z \in \mathbb{H}} \{\langle y, z \rangle - g(z)\},$$

138 for all $y \in \mathbb{H}$. Recall that g^* is closed and convex.

We denote by $\Gamma_0(\mathbb{H})$ the set of all extended-real-valued functions $g : \mathbb{H} \rightarrow \mathbb{R} \cup \{+\infty\}$ that are proper closed and convex. If $g \in \Gamma_0(\mathbb{H})$, recall that $g^* \in \Gamma_0(\mathbb{H})$. The *Fenchel-Moreau equality* is given by $g^{**} = g$. For all $g \in \Gamma_0(\mathbb{H})$, we denote by $\partial g : \mathbb{H} \rightrightarrows \mathbb{H}$ the *Fenchel-Moreau subdifferential* of g defined by

$$\partial g(x) := \{y \in \mathbb{H} \mid \langle y, z - x \rangle \leq g(z) - g(x), \forall z \in \mathbb{H}\},$$

139 for all $x \in \mathbb{H}$. It is easy to check that ∂g is a monotone operator and that, for
 140 all $x \in \mathbb{H}$, $0 \in \partial g(x)$ if and only if $x \in \text{argmin } g$. Moreover, for all $x, y \in \mathbb{H}$, it holds
 141 that $y \in \partial g(x)$ if and only if $x \in \partial g^*(y)$. Recall that, if g is differentiable on \mathbb{H} ,
 142 then $\partial g(x) = \{\nabla g(x)\}$ for all $x \in \mathbb{H}$.

143 Let $A : \mathbb{H} \rightarrow \mathbb{H}$ be a single-valued operator defined everywhere on \mathbb{H} , and let $g \in \Gamma_0(\mathbb{H})$.
 144 We denote by $\text{VI}(A, g)$ the variational inequality which consists of finding $y \in \mathbb{H}$ such
 145 that

$$146 \quad -A(y) \in \partial g(y),$$

147 or equivalently,

$$148 \quad \langle A(y), z - y \rangle + g(z) - g(y) \geq 0,$$

149 for all $z \in \mathbb{H}$. Then we denote by $\text{Sol}_{\text{VI}}(A, g)$ the set of solutions of $\text{VI}(A, g)$. Re-
 150 call that if A is Lipschitzian and strongly monotone, then $\text{VI}(A, g)$ admits a unique
 151 solution, *i.e.* $\text{Sol}_{\text{VI}}(A, g)$ is a singleton.

Let $g \in \Gamma_0(\mathbb{H})$. The classical *proximity operator* of g is defined by

$$\text{prox}_g := (\text{I} + \partial g)^{-1}.$$

Recall that prox_g is a single-valued operator defined everywhere on \mathbb{H} . Moreover, it can be characterized as follows:

$$\text{prox}_g(x) = \text{argmin} \left(g + \frac{1}{2} \|\cdot - x\|^2 \right) = \text{Sol}_{\text{VI}}(-L_x, g),$$

for all $x \in H$. It is also well-known that

$$\text{Fix}(\text{prox}_g) = \text{argmin } g.$$

The classical *Moreau's envelope* $M_g : H \rightarrow \mathbb{R}$ of g is defined by

$$M_g(x) := \min \left(g + \frac{1}{2} \| \cdot - x \|^2 \right),$$

for all $x \in H$. Recall that M_g is convex and differentiable on H with $\nabla M_g = \text{prox}_{g^*}$. Let us also recall the classical Moreau's decompositions

$$\text{prox}_g + \text{prox}_{g^*} = \text{I} \quad \text{and} \quad M_g + M_{g^*} = \frac{1}{2} \| \cdot \|^2.$$

152 Finally, it is well-known that if $g = \iota_K$ is the *indicator function* of a nonempty closed
153 and convex subset K of H , that is, $\iota_K(x) = 0$ if $x \in K$ and $\iota_K(x) = +\infty$ if not,
154 then $\text{prox}_g = \text{proj}_K$, where proj_K denotes the classical projection operator on K .

155 2. The f -proximity operator.

156 **2.1. Main result.** Let $f, g \in \Gamma_0(H)$. In this section we introduce (see Def-
157 inition 2.1) a new operator denoted by prox_g^f , generalizing the classical proximity
158 operator prox_g . Under the additivity condition $\partial(f + g) = \partial f + \partial g$, we prove in
159 Theorem 2.7 that prox_{f+g} can be written as the composition of prox_f with prox_g^f .

DEFINITION 2.1 (f -proximity operator). *Let $f, g \in \Gamma_0(H)$. The f -proximity operator of g is the set-valued map $\text{prox}_g^f : H \rightrightarrows H$ defined by*

$$\text{prox}_g^f := (\text{I} + \partial g \circ \text{prox}_f)^{-1}.$$

160 Note that prox_g^f can be seen as a generalization of prox_g since $\text{prox}_g^c = \text{prox}_g$ for all
161 constant $c \in \mathbb{R}$.

162 *Example 2.2.* Let us assume that $H = \mathbb{R}$. We consider $f = \iota_{[-1,1]}$ and $g(x) = |x|$ for
163 all $x \in \mathbb{R}$. In that case we obtain that $\partial g \circ \text{prox}_f = \partial g$ and thus $\text{prox}_g^f = \text{prox}_g$.

164 Example 2.2 provides a simple situation where $\text{prox}_g^f = \text{prox}_g$ while f is not constant.
165 We provide in Propositions 2.11 and 2.14 some general sufficient (and necessary)
166 conditions under which $\text{prox}_g^f = \text{prox}_g$.

167 *Example 2.3.* Let us assume that $H = \mathbb{R}$. We consider $f = \iota_{\{0\}}$ and $g(x) = |x|$ for
168 all $x \in \mathbb{R}$. In that case we obtain that $\partial g \circ \text{prox}_f(x) = [-1, 1]$ for all $x \in \mathbb{R}$. As
169 a consequence $\text{prox}_g^f(x) = [x - 1, x + 1]$ for all $x \in \mathbb{R}$. See Figure 1 for graphic
170 representations of prox_g and prox_g^f in that case.

171 Example 2.3 provides a simple illustration where prox_g^f is not single-valued. In
172 particular it follows that prox_g^f cannot be written as a proximity operator prox_φ
173 with $\varphi \in \Gamma_0(H)$. We provide in Proposition 2.17 some sufficient conditions under
174 which prox_g^f is single-valued. Moreover, Example 2.3 provides a simple situation
175 where $\partial g \circ \text{prox}_f$ is not a monotone operator. As a consequence, it may be possible
176 that $D(\text{prox}_g^f) \subsetneq H$. In the next proposition, a necessary and sufficient condition
177 under which $D(\text{prox}_g^f) = H$ is derived.

178 PROPOSITION 2.4. *Let $f, g \in \Gamma_0(H)$. It holds that $D(\text{prox}_g^f) = H$ if and only if the*
179 *additivity condition*

$$180 \quad (1) \quad \partial(f + g) = \partial f + \partial g,$$

181 *is satisfied.*

182 *Proof.* We first assume that $\partial(f + g) = \partial f + \partial g$. Let $x \in H$. Defining $w =$
 183 $\text{prox}_{f+g}(x) \in H$, we obtain that $x \in w + \partial(f + g)(w) = w + \partial f(w) + \partial g(w)$. Thus,
 184 there exist $w_f \in \partial f(w)$ and $w_g \in \partial g(w)$ such that $x = w + w_f + w_g$. We de-
 185 fine $y = w + w_f \in w + \partial f(w)$. In particular we have $w = \text{prox}_f(y)$. Moreover we
 186 obtain $x = y + w_g \in y + \partial g(w) = y + \partial g(\text{prox}_f(y))$. We conclude that $y \in \text{prox}_g^f(x)$.

187 Without any additional assumption and directly from the definition of the subdiffer-
 188 ential, one can easily see that the inclusion $\partial f(w) + \partial g(w) \subset \partial(f + g)(w)$ is always
 189 satisfied for every $w \in H$. Now let us assume that $D(\text{prox}_g^f) = H$. Let $w \in H$ and
 190 let $z \in \partial(f + g)(w)$. We consider $x = w + z \in w + \partial(f + g)(w)$. In particular it holds
 191 that $w = \text{prox}_{f+g}(x)$. Since $D(\text{prox}_g^f) = H$, there exists $y \in \text{prox}_g^f(x)$ and thus it
 192 holds that $x \in y + \partial g(\text{prox}_f(y))$. Moreover, since $y \in \text{prox}_f(y) + \partial f(\text{prox}_f(y))$, we get
 193 that $x \in \text{prox}_f(y) + \partial f(\text{prox}_f(y)) + \partial g(\text{prox}_f(y)) \subset \text{prox}_f(y) + \partial(f + g)(\text{prox}_f(y))$.
 194 Thus it holds that $\text{prox}_f(y) = \text{prox}_{f+g}(x) = w$. Moreover, since $x \in \text{prox}_f(y) +$
 195 $\partial f(\text{prox}_f(y)) + \partial g(\text{prox}_f(y))$, we obtain that $x \in w + \partial f(w) + \partial g(w)$. We have proved
 196 that $z = x - w \in \partial f(w) + \partial g(w)$. This concludes the proof. \square

197 In most of the present paper, we will assume that Condition (1) is satisfied. It is not
 198 our aim here to discuss the weakest qualification condition ensuring Condition (1). A
 199 wide literature already deals with this topic (see, e.g., [1, 11, 20]). However, we recall
 200 in the following remark the classical sufficient condition of Moreau-Rockafellar under
 201 which Condition (1) holds true (see, e.g., [2, Corollary 16.38 p.234]), and we provide
 202 a simple example where Condition (1) does not hold and $D(\text{prox}_g^f) \subsetneq H$.

203 *Remark 2.5* (Moreau-Rockafellar theorem). Let $f, g \in \Gamma_0(H)$ such that $\text{dom}(f) \cap$
 204 $\text{int}(\text{dom}(g)) \neq \emptyset$. Then $\partial(f + g) = \partial f + \partial g$.

205 *Example 2.6.* Let us assume that $H = \mathbb{R}$. We consider $f = \iota_{\mathbb{R}^-}$ and $g(x) = \iota_{\mathbb{R}^+}(x) - \sqrt{x}$
 206 for all $x \in \mathbb{R}$. In that case, one can easily check that $\partial f(0) + \partial g(0) = \emptyset \subsetneq \mathbb{R} =$
 207 $\partial(f + g)(0)$ and $D(\text{prox}_g^f) = \emptyset \subsetneq H$.

208 We are now in position to state and prove the main result of the present paper.

THEOREM 2.7. *Let $f, g \in \Gamma_0(H)$ such that $\partial(f + g) = \partial f + \partial g$. It holds that*

$$\text{prox}_{f+g} = \text{prox}_f \circ \text{prox}_g^f.$$

209 *In other words, for every $x \in H$, we have $\text{prox}_{f+g}(x) = \text{prox}_f(z)$ for all $z \in \text{prox}_g^f(x)$.*

Proof. Let $x \in H$ and let $y \in \text{prox}_g^f(x)$ constructed as in the first part of the proof of
 Proposition 2.4. In particular it holds that $\text{prox}_f(y) = \text{prox}_{f+g}(x)$. Let $z \in \text{prox}_g^f(x)$.
 We know that $x - y \in \partial g(\text{prox}_f(y))$ and $x - z \in \partial g(\text{prox}_f(z))$. Since ∂g is a monotone
 operator, we obtain that

$$\langle (x - y) - (x - z), \text{prox}_f(y) - \text{prox}_f(z) \rangle \geq 0.$$

From the cocoercivity (see [2, Definition 4.4 p.60]) of the proximity operator, we
 obtain that

$$0 \geq \langle y - z, \text{prox}_f(y) - \text{prox}_f(z) \rangle \geq \|\text{prox}_f(y) - \text{prox}_f(z)\|^2 \geq 0.$$

210 We deduce that $\text{prox}_f(z) = \text{prox}_f(y) = \text{prox}_{f+g}(x)$. The proof is complete. \square

211 *Remark 2.8.* Let $f, g \in \Gamma_0(H)$ such that $\partial(f + g) = \partial f + \partial g$ and let $x \in H$. The-
 212 orem 2.7 states that, even if $\text{prox}_g^f(x)$ is not a singleton, all elements of $\text{prox}_g^f(x)$

213 has the same value through the proximity operator prox_f , and this value is equal
214 to $\text{prox}_{f+g}(x)$.

215 *Remark 2.9.* Note that the additivity condition $\partial(f+g) = \partial f + \partial g$ is not only suf-
216 ficient, but also necessary for the validity of the equality $\text{prox}_{f+g}^f = \text{prox}_f \circ \text{prox}_g^f$.
217 Indeed, from Proposition 2.4, if $\partial f + \partial g \subsetneq \partial(f+g)$, then there exists $x \in \mathbf{H}$ such that
218 $\text{prox}_g^f(x) = \emptyset$ and thus $\text{prox}_{f+g}(x) \neq \text{prox}_f \circ \text{prox}_g^f(x)$.

219 *Remark 2.10.* Let $f, g \in \Gamma_0(\mathbf{H})$ such that $\partial(f+g) = \partial f + \partial g$. From Theorem 2.7, we
220 deduce that $\mathbf{R}(\text{prox}_{f+g}) \subset \mathbf{R}(\text{prox}_f) \cap \mathbf{R}(\text{prox}_g)$. If the additivity condition $\partial(f+g) =$
221 $\partial f + \partial g$ is not satisfied, this remark does not hold true anymore. Indeed, with the
222 framework of Example 2.6, we have $\mathbf{R}(\text{prox}_{f+g}) = \{0\}$ while $0 \notin \mathbf{R}(\text{prox}_g)$.

223 **2.2. Additional results.** Let $f, g \in \Gamma_0(\mathbf{H})$. We know that prox_g^f is a gener-
224 alization of prox_g in the sense that $\text{prox}_g^f = \text{prox}_g$ if f is constant for instance. In
225 the next proposition, our aim is to provide more general sufficient (and necessary)
226 conditions under which $\text{prox}_g^f = \text{prox}_g$. We will base our discussion on the following
227 conditions:

$$228 \quad (2) \quad \forall x \in \mathbf{H}, \quad \partial g(x) \subset \partial g(\text{prox}_f(x)),$$

229

$$230 \quad (3) \quad \forall x \in \mathbf{H}, \quad \partial g(\text{prox}_f(x)) \subset \partial g(x).$$

231 Note that Condition (2) has been introduced by Y.-L. Yu in [24] as a sufficient con-
232 dition under which $\text{prox}_{f+g} = \text{prox}_f \circ \text{prox}_g$.

233 **PROPOSITION 2.11.** *Let $f, g \in \Gamma_0(\mathbf{H})$.*

234 (i) *If Condition (2) is satisfied, then $\text{prox}_g(x) \in \text{prox}_g^f(x)$ for all $x \in \mathbf{H}$.*

235 (ii) *If Conditions (1) and (3) are satisfied, then $\text{prox}_g^f(x) = \text{prox}_g(x)$ for all*
236 *$x \in \mathbf{H}$.*

237 *In both cases, it holds that $\text{prox}_{f+g} = \text{prox}_f \circ \text{prox}_g$.*

238 *Proof.* Let $x \in \mathbf{H}$. If Condition (2) is satisfied, considering $y = \text{prox}_g(x)$, we get
239 that $x \in y + \partial g(y) \subset y + \partial g(\text{prox}_f(y))$ and thus $y \in \text{prox}_g^f(x)$. In particular, it holds
240 that $\mathbf{D}(\text{prox}_g^f) = \mathbf{H}$ and thus $\partial(f+g) = \partial f + \partial g$ from Proposition 2.4. On the other
241 hand, if Conditions (1) and (3) are satisfied, then $\mathbf{D}(\text{prox}_g^f) = \mathbf{H}$ from Proposition 2.4.
242 Considering $y \in \text{prox}_g^f(x)$, we get that $x \in y + \partial g(\text{prox}_f(y)) \subset y + \partial g(y)$ and thus $y =$
243 $\text{prox}_g(x)$. The last assertion of Proposition 2.11 directly follows from Theorem 2.7. \square

244 *Remark 2.12.* From Proposition 2.11, we deduce that Condition (2) implies Condi-
245 tion (1).

246 In the first item of Proposition 2.11 and if prox_g^f is set-valued, we are in the situation
247 where prox_g is a selection of prox_g^f . Proposition 2.14 specifies this selection in the
248 case where $\partial(f+g) = \partial f + \partial g$.

249 **LEMMA 2.13.** *Let $f, g \in \Gamma_0(\mathbf{H})$. Then $\text{prox}_g^f(x)$ is a nonempty closed and convex*
250 *subset of \mathbf{H} for all $x \in \mathbf{D}(\text{prox}_g^f)$.*

251 *Proof.* The proof of Lemma 2.13 is provided after the proof of Proposition 3.2 (re-
252 quired). \square

PROPOSITION 2.14. *Let $f, g \in \Gamma_0(\mathbf{H})$ such that $\partial(f+g) = \partial f + \partial g$ and let $x \in \mathbf{H}$.
If $\text{prox}_g(x) \in \text{prox}_g^f(x)$, then*

$$\text{prox}_g(x) = \text{proj}_{\text{prox}_g^f(x)}(\text{prox}_{f+g}(x)).$$

253 *Proof.* If $\text{prox}_g(x) \in \text{prox}_g^f(x)$, then $x \in \text{D}(\text{prox}_g^f)$ and thus $\text{prox}_g^f(x)$ is a nonempty
 254 closed and convex subset of H from Lemma 2.13. Let $z \in \text{prox}_g^f(x)$. In par-
 255 ticular we have $\text{prox}_f(z) = \text{prox}_{f+g}(x)$ from Theorem 2.7. Using the fact that
 256 $x - \text{prox}_g(x) \in \partial g(\text{prox}_g(x))$ and $x - z \in \partial g(\text{prox}_f(z)) = \partial g(\text{prox}_{f+g}(x))$ together
 257 with the monotonicity of ∂g , we obtain that

$$\begin{aligned}
 258 \quad & \langle \text{prox}_{f+g}(x) - \text{prox}_g(x), z - \text{prox}_g(x) \rangle \\
 259 \quad & = \langle \text{prox}_{f+g}(x) - \text{prox}_g(x), (x - \text{prox}_g(x)) - (x - z) \rangle \leq 0.
 \end{aligned}$$

262 Since $\text{prox}_g(x) \in \text{prox}_g^f(x)$, we conclude the proof from the classical characterization
 263 of $\text{proj}_{\text{prox}_g^f(x)}$. \square

Remark 2.15. Let $f = \iota_{\{\omega\}}$ with $\omega \in \text{H}$ and let $g \in \Gamma_0(\text{H})$ such that $\omega \in \text{int}(\text{dom}(g))$.
 Hence the additivity condition $\partial(f+g) = \partial f + \partial g$ is satisfied from Remark 2.5. From
 Remark 2.10 and since $\text{prox}_f = \text{proj}_{\{\omega\}}$, we easily deduce that $\text{R}(\text{prox}_{f+g}) = \{\omega\}$.
 Let $x \in \text{H}$ such that $\text{prox}_g(x) \in \text{prox}_g^f(x)$. From Proposition 2.14 we get that

$$\text{prox}_g(x) = \text{proj}_{\text{prox}_g^f(x)}(\omega).$$

264 If moreover $\omega = 0$, we deduce that $\text{prox}_g(x)$ is the particular selection that corresponds
 265 to the element of minimal norm in $\text{prox}_g^f(x)$ (also known as the *lazy selection*). The
 266 following example is in this sense.

267 *Example 2.16.* Let us consider the framework of Example 2.3. In that case, Condi-
 268 tions (1) and (2) are satisfied. We deduce from Proposition 2.11 that $\text{prox}_g(x) \in$
 269 $\text{prox}_g^f(x)$ for all $x \in \mathbb{R}$. From Remark 2.15, we conclude that $\text{prox}_g(x)$ is exactly the
 270 element of minimal norm in $\text{prox}_g^f(x)$ for all $x \in \mathbb{R}$. This result is clearly illustrated
 by the graphs of prox_g and prox_g^f provided in Figure 1.

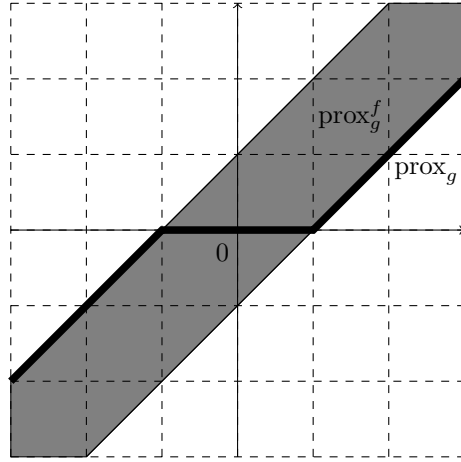


FIG. 1. Examples 2.3 and 2.16, graph of prox_g in bold line, and graph of prox_g^f in gray.

271
 272 Let $f, g \in \Gamma_0(\text{H})$ such that $\partial(f+g) = \partial f + \partial g$. From Theorem 2.7, one can easily
 273 see that, if prox_f is injective, then prox_g^f is single-valued. Since the injection of prox_f
 274 is too restrictive, other sufficient conditions under which prox_g^f is single-valued are
 275 provided from Theorem 2.7 in the next proposition.

276 PROPOSITION 2.17. *Let $f, g \in \Gamma_0(\mathbb{H})$ such that $\partial(f + g) = \partial f + \partial g$. If either ∂f
277 or ∂g is single-valued, then prox_g^f is single-valued.*

278 *Proof.* Let $x \in \mathbb{H}$ and let $z_1, z_2 \in \text{prox}_g^f(x)$. From Theorem 2.7, it holds that
279 $\text{prox}_f(z_1) = \text{prox}_f(z_2) = \text{prox}_{f+g}(x)$. If the operator ∂f is single-valued, we ob-
280 tain that $z_1 = \text{prox}_{f+g}(x) + \partial f(\text{prox}_{f+g}(x)) = z_2$. If the operator ∂g is single-valued,
281 we get $x - z_1 = \partial g(\text{prox}_f(z_1)) = \partial g(\text{prox}_f(z_2)) = x - z_2$ and thus $z_1 = z_2$. \square

3. Relations with the Douglas-Rachford operator. Let $f, g \in \Gamma_0(\mathbb{H})$. The classical *Douglas-Rachford operator* $\mathcal{T}_{f,g} : \mathbb{H} \rightarrow \mathbb{H}$ associated to f and g is usually defined by

$$\mathcal{T}_{f,g}(y) := y - \text{prox}_f(y) + \text{prox}_g(2\text{prox}_f(y) - y),$$

282 for all $y \in \mathbb{H}$. We refer to [9, 10, 15] and to [2, Section 27.2 p.400] for more details.

One aim of this section is to study the relations between the f -proximity operator prox_g^f introduced in this paper and the Douglas-Rachford operator $\mathcal{T}_{f,g}$. For this purpose, we introduce an extension $\overline{\mathcal{T}}_{f,g} : \mathbb{H} \times \mathbb{H} \rightarrow \mathbb{H}$ of the classical Douglas-Rachford operator defined by

$$\overline{\mathcal{T}}_{f,g}(x, y) := y - \text{prox}_f(y) + \text{prox}_g(x + \text{prox}_f(y) - y),$$

283 for all $x, y \in \mathbb{H}$.

284 Note that $\mathcal{T}_{f,g}(y) = \overline{\mathcal{T}}_{f,g}(\text{prox}_f(y), y)$ for all $y \in \mathbb{H}$, and that the definition of $\overline{\mathcal{T}}_{f,g}$
285 only depends on the knowledge of prox_f and prox_g .

286 **3.1. Several characterizations of prox_g^f .** Let $f, g \in \Gamma_0(\mathbb{H})$. In this section,
287 our aim is to derive several characterizations of prox_g^f in terms of solutions of varia-
288 tional inequalities, of minimization problems and of fixed point problems (see Propo-
289 sition 3.2).

LEMMA 3.1. *Let $f, g \in \Gamma_0(\mathbb{H})$. It holds that*

$$\overline{\mathcal{T}}_{f,g}(x, \cdot) = \text{prox}_{g^* \circ L_x} \circ \text{prox}_{f^*},$$

290 for all $x \in \mathbb{H}$.

291 *Proof.* Let $x \in \mathbb{H}$. Lemma 3.1 directly follows from the equality $\text{prox}_{g^* \circ L_x} = L_x \circ$
292 $\text{prox}_{g^*} \circ L_x$ (see [2, Proposition 23.29 p.342]) and from Moreau's decompositions. \square

PROPOSITION 3.2. *Let $f, g \in \Gamma_0(\mathbb{H})$. It holds that*

$$\text{prox}_g^f(x) = \text{Sol}_{\text{VI}}(\text{prox}_f, g^* \circ L_x) = \text{argmin}(\text{M}_{f^*} + g^* \circ L_x) = \text{Fix}(\overline{\mathcal{T}}_{f,g}(x, \cdot)),$$

293 for all $x \in \mathbb{H}$.

294 *Proof.* In this proof we will use standard properties of convex analysis recalled in
295 Section 1.2. Let $x \in \mathbb{H}$. One can easily prove that $\partial(g^* \circ L_x) = -\partial g^* \circ L_x$. For
296 all $y \in \mathbb{H}$, it holds that

$$\begin{aligned} 297 \quad y \in \text{prox}_g^f(x) &\iff x - y \in \partial g(\text{prox}_f(y)) \\ 298 &\iff \text{prox}_f(y) \in \partial g^*(x - y) \\ 299 &\iff -\text{prox}_f(y) \in \partial(g^* \circ L_x)(y). \end{aligned}$$

300 Moreover, since $\text{dom}(\text{M}_{f^*}) = \mathbb{H}$ and from Remark 2.5, we have

$$\begin{aligned} 301 \quad -\text{prox}_f(y) \in \partial(g^* \circ L_x)(y) &\iff 0 \in \nabla \text{M}_{f^*}(y) + \partial(g^* \circ L_x)(y) \\ 302 &\iff 0 \in \partial(\text{M}_{f^*} + g^* \circ L_x)(y). \end{aligned}$$

303 Finally,

$$\begin{aligned}
 304 \quad & -\text{prox}_f(y) \in \partial(g^* \circ L_x)(y) \iff \text{prox}_{f^*}(y) \in y + \partial(g^* \circ L_x)(y) \\
 305 \quad & \iff y = \text{prox}_{g^* \circ L_x} \circ \text{prox}_{f^*}(y).
 \end{aligned}$$

306 This concludes the proof from Lemma 3.1. \square

Proof of Lemma 2.13. Let $x \in \text{D}(\text{prox}_g^f)$. In particular $\text{prox}_g^f(x)$ is not empty. From Proposition 3.2, we have

$$\text{prox}_g^f(x) = \text{argmin}(\text{M}_{f^*} + g^* \circ L_x).$$

307 Since $\text{M}_{f^*} + g^* \circ L_x \in \Gamma_0(\text{H})$, one can easily deduce that $\text{prox}_g^f(x)$ is closed and
 308 convex. \square

309 **3.2. A one-loop algorithm in order to compute prox_g^f numerically.** Let
 310 $f, g \in \Gamma_0(\text{H})$. In this section, our aim is to derive a one-loop algorithm, that depends
 311 only on the knowledge of prox_f and prox_g , allowing to compute numerically an element
 312 of $\text{prox}_g^f(x)$ for all $x \in \text{D}(\text{prox}_g^f)$. We refer to Algorithm (\mathcal{A}_1) in Theorem 3.3.

313 Moreover, if the additivity condition $\partial(f + g) = \partial f + \partial g$ is satisfied, it follows from
 314 Theorem 2.7 that Algorithm (\mathcal{A}_1) is a one-loop algorithm allowing to compute numerically
 315 $\text{prox}_{f+g}(x)$ for all $x \in \text{H}$ with the only knowledge of prox_f and prox_g .

316 **THEOREM 3.3.** *Let $f, g \in \Gamma_0(\text{H})$ and let $x \in \text{D}(\text{prox}_g^f)$ be fixed. Then, Algorithm (\mathcal{A}_1)*
 317 *given by*

$$318 \quad (\mathcal{A}_1) \quad \begin{cases} y_0 \in \text{H}, \\ y_{k+1} = \overline{\mathcal{T}}_{f,g}(x, y_k), \end{cases}$$

319 *weakly converges to an element $y^* \in \text{prox}_g^f(x)$. Moreover, if the additivity condition*
 320 *$\partial(f + g) = \partial f + \partial g$ is satisfied, it holds that $\text{prox}_f(y^*) = \text{prox}_{f+g}(x)$.*

321 *Proof.* From Lemma 3.1, $\overline{\mathcal{T}}_{f,g}(x, \cdot)$ coincides with the composition of two firmly non-
 322 expansive operators, and thus of two non-expansive and $\frac{1}{2}$ -averaged operators (see [2,
 323 Remark 4.24(iii) p.68]). Since $x \in \text{D}(\text{prox}_g^f)$, it follows from Proposition 3.2 and
 324 Lemma 3.1 that $\text{Fix}(\text{prox}_{g^* \circ L_x} \circ \text{prox}_{f^*}) \neq \emptyset$. We conclude from [2, Theorem 5.22
 325 p.82] that Algorithm (\mathcal{A}_1) weakly converges to a fixed point y^* of $\overline{\mathcal{T}}_{f,g}(x, \cdot)$. From
 326 Proposition 3.2, it holds that $y^* \in \text{prox}_g^f(x)$. Finally, if the additivity condition
 327 $\partial(f + g) = \partial f + \partial g$ is satisfied, we conclude that $\text{prox}_f(y^*) = \text{prox}_{f+g}(x)$ from
 328 Theorem 2.7. \square

Remark 3.4. Let $f, g \in \Gamma_0(\text{H})$ and let $x \in \text{D}(\text{prox}_g^f)$. Algorithm (\mathcal{A}_1) consists in a
 fixed-point algorithm from the characterization given in Proposition 3.2 by

$$\text{prox}_g^f(x) = \text{Fix}(\overline{\mathcal{T}}_{f,g}(x, \cdot)).$$

Actually, one can easily see that Algorithm (\mathcal{A}_1) also coincides with the well-known
Forward-Backward algorithm (see [7, Section 10.3 p.191] for details) from the charac-
 terization given in Proposition 3.2 by

$$\text{prox}_g^f(x) = \text{argmin}(\text{M}_{f^*} + g^* \circ L_x).$$

329 Indeed, we recall that M_{f^*} is differentiable with $\nabla \text{M}_{f^*} = \text{prox}_f$. We also refer to
 330 Section 4.1 for a brief discussion about the Forward-Backward algorithm.

331 Let $f, g \in \Gamma_0(\mathbb{H})$. As mentioned in the introduction, no proximal algorithm $x_{n+1} =$
 332 $\text{prox}_{f+g}(x_n)$, using only the knowledge of prox_f and prox_g , has been provided in the
 333 literature. This remains a very interesting open challenge in the literature. However,
 334 we will introduce now the notion of *proximal-like algorithm* (see Definition 3.5) and
 335 we will provide in Remark 3.7 such a proximal-like algorithm $x_{n+1} = \text{prox}_{f+g}(x_n)$
 336 requiring only the knowledge of prox_f and prox_g .

337 **DEFINITION 3.5** (Proximal-like algorithm). *Let $g \in \Gamma_0(\mathbb{H})$. An algorithm is said to*
 338 *be a proximal-like algorithm $x_{n+1} = \text{prox}_g(x_n)$ if it can be written as*

$$339 \quad \left\{ \begin{array}{l} x_0 \in \mathbb{H}, \\ x_{n+1} = P_1(y_n^*), \\ \quad \text{where } y_n^* \text{ is given by solving a weakly} \\ \quad \text{convergent auxiliary subalgorithm} \\ \left\{ \begin{array}{l} y_{n,0} \in \mathbb{H}, \\ y_{n,k+1} = P_2(x_n, y_{n,k}), \end{array} \right. \end{array} \right.$$

where $P_1 : \mathbb{H} \rightarrow \mathbb{H}$ and $P_2 : \mathbb{H} \times \mathbb{H} \rightarrow \mathbb{H}$ are two given operators satisfying

$$P_1(\text{Fix}(P_2(x, \cdot))) = \text{prox}_g(x),$$

340 for all $x \in \mathbb{H}$.

341 **Remark 3.6.** In contrary to the classical proximal, Douglas-Rachford and Forward-
 342 Backward algorithms, it should be noted that a proximal-like algorithm is a two-loops
 343 algorithm.

344 **Remark 3.7.** Let $f, g \in \Gamma_0(\mathbb{H})$ such that $\partial(f + g) = \partial f + \partial g$. From Theorem 2.7,
 345 Proposition 3.2 and Theorem 3.3, Algorithm (\mathcal{A}_2) given by

$$346 \quad (\mathcal{A}_2) \quad \left\{ \begin{array}{l} x_0 \in \mathbb{H}, \\ x_{n+1} = \text{prox}_f(y_n^*), \\ \quad \text{where } y_n^* \text{ is given by solving the weakly} \\ \quad \text{convergent auxiliary subalgorithm} \\ \left\{ \begin{array}{l} y_{n,0} \in \mathbb{H}, \\ y_{n,k+1} = \overline{\mathcal{T}}_{f,g}(x_n, y_{n,k}), \end{array} \right. \end{array} \right.$$

347 is a *proximal-like algorithm* $x_{n+1} = \text{prox}_{f+g}(x_n)$ that only requires the knowledge
 348 of prox_f and prox_g .

349 **Remark 3.8.** As mentioned in the introduction, the aim of the present theoretical pa-
 350 per is not to discuss numerical experiments and comparisons between numerical algo-
 351 rithms (this should be the topic of future works). However, in contrary to the classical
 352 Douglas-Rachford algorithm, it should be noted that Algorithm (\mathcal{A}_2) is a two-loops
 353 algorithm. As a consequence, it should not be expected from Algorithm (\mathcal{A}_2) bet-
 354 ter performances than the Douglas-Rachford algorithm for solving the minimization
 355 problem $\text{argmin } f + g$.

3.3. An additional result on the Douglas-Rachford operator. Let $f, g \in \Gamma_0(\mathbb{H})$. It is well-known in the literature (and it can be easily proved) that

$$\text{prox}_f(\text{Fix}(\mathcal{T}_{f,g})) \subset \text{argmin } f + g.$$

356 Our aim in this section is to prove, with the help of Theorem 2.7, that the opposite
 357 inclusion holds true under the additivity condition $\partial(f + g) = \partial f + \partial g$. To the best
 358 of our knowledge, this result is new in the literature.

LEMMA 3.9. *Let $f, g \in \Gamma_0(\mathbb{H})$. It holds that*

$$\text{Fix}(\mathcal{T}_{f,g}) = \text{Fix}(\text{prox}_g^f \circ \text{prox}_f).$$

359 *Proof.* Let $z \in \mathbb{H}$. It holds from Proposition 3.2 that

$$\begin{aligned} 360 \quad z \in \text{Fix}(\mathcal{T}_{f,g}) &\iff z = \mathcal{T}_{f,g}(z) = \overline{\mathcal{T}}_{f,g}(\text{prox}_f(z), z) \\ 361 &\iff z \in \text{Fix}(\overline{\mathcal{T}}_{f,g}(\text{prox}_f(z), \cdot)) = \text{prox}_g^f(\text{prox}_f(z)) \\ 362 &\iff z \in \text{Fix}(\text{prox}_g^f \circ \text{prox}_f). \end{aligned}$$

363 The proof is complete. □

PROPOSITION 3.10. *Let $f, g \in \Gamma_0(\mathbb{H})$ such that $\partial(f + g) = \partial f + \partial g$. It holds that*

$$\text{argmin } f + g = \text{prox}_f(\text{Fix}(\mathcal{T}_{f,g})).$$

364 *Proof.* Let $y \in \text{Fix}(\mathcal{T}_{f,g})$. Then $y \in \text{Fix}(\text{prox}_g^f \circ \text{prox}_f)$ from Lemma 3.9. Thus
 365 $y \in \text{prox}_g^f \circ \text{prox}_f(y)$. From Theorem 2.7, we get that $\text{prox}_f(y) = \text{prox}_{f+g}(\text{prox}_f(y))$
 366 and thus $\text{prox}_f(y) \in \text{argmin } f + g$.

367 Let $x \in \text{argmin } f + g$. Since $D(\text{prox}_g^f) = \mathbb{H}$ from Proposition 2.4, let us consider
 368 $y \in \text{prox}_g^f(x)$. From Theorem 2.7, it holds that $x = \text{prox}_{f+g}(x) = \text{prox}_f(y)$. Let
 369 us prove that $y \in \text{Fix}(\mathcal{T}_{f,g})$. Since $y \in \text{prox}_g^f(x) = \text{Fix}(\overline{\mathcal{T}}_{f,g}(x, \cdot))$, we get that
 370 $y = \overline{\mathcal{T}}_{f,g}(x, y) = \overline{\mathcal{T}}_{f,g}(\text{prox}_f(y), y) = \mathcal{T}_{f,g}(y)$. The proof is complete. □

371 **4. Some other applications and forthcoming works.** This section can be
 372 seen as a conclusion of the paper. Its aim is to provide a glimpse of some other appli-
 373 cations of our main result (Theorem 2.7) and to raise open questions for forthcoming
 374 works. This section is splitted into two parts.

375 **4.1. Relations with the classical Forward-Backward operator.** Let $f,$
 376 $g \in \Gamma_0(\mathbb{H})$ such that g is differentiable on \mathbb{H} . In that situation, note that the additivity
 377 condition $\partial(f + g) = \partial f + \partial g$ is satisfied from Remark 2.5, and that Proposition 2.17
 378 implies that prox_g^f is single-valued.

In that framework, the classical *Forward-Backward operator* $\mathcal{F}_{f,g} : \mathbb{H} \rightarrow \mathbb{H}$ associated
 to f and g is usually defined by

$$\mathcal{F}_{f,g}(y) := \text{prox}_f(y - \nabla g(y)),$$

for all $y \in \mathbb{H}$. We refer to [7, Section 10.3 p.191] for more details. Let us introduce
 the extension $\overline{\mathcal{F}}_{f,g} : \mathbb{H} \times \mathbb{H} \rightarrow \mathbb{H}$ defined by

$$\overline{\mathcal{F}}_{f,g}(x, y) := \text{prox}_f(x - \nabla g(y)),$$

379 for all $x, y \in \mathbb{H}$. In particular, it holds that $\mathcal{F}_{f,g}(y) = \overline{\mathcal{F}}_{f,g}(y, y)$ for all $y \in \mathbb{H}$. The
 380 following result follows from Theorem 2.7.

PROPOSITION 4.1. *Let $f, g \in \Gamma_0(\mathbb{H})$ such that g is differentiable on \mathbb{H} . Then*

$$\text{prox}_{f+g}(x) = \text{Fix}(\overline{\mathcal{F}}_{f,g}(x, \cdot)),$$

381 *for all $x \in \mathbb{H}$.*

382 *Proof.* Let $x \in \mathbb{H}$. Firstly, let $z = \text{prox}_{f+g}(x)$ and let $y = \text{prox}_g^f(x)$. In particular, we
 383 have $x = y + \nabla g(\text{prox}_f(y))$. From Theorem 2.7, we get that $z = \text{prox}_f(y) = \text{prox}_f(x -$
 384 $\nabla g(\text{prox}_f(y))) = \text{prox}_f(x - \nabla g(z)) = \overline{\mathcal{F}}_{f,g}(x, z)$. Conversely, let $z \in \text{Fix}(\overline{\mathcal{F}}_{f,g}(x, \cdot))$,
 385 that is, $z = \text{prox}_f(x - \nabla g(z))$. Considering $y = x - \nabla g(z)$, we have $z = \text{prox}_f(y)$ and
 386 thus $x = y + \nabla g(\text{prox}_f(y))$, that is, $y = \text{prox}_g^f(x)$. Finally, from Theorem 2.7, we get
 387 that $z = \text{prox}_f \circ \text{prox}_g^f(x) = \text{prox}_{f+g}(x)$. \square

388 From Proposition 4.1, we retrieve the following classical result.

PROPOSITION 4.2. *Let $f, g \in \Gamma_0(\mathbb{H})$ such that g is differentiable on \mathbb{H} . Then*

$$\text{argmin } f + g = \text{Fix}(\mathcal{F}_{f,g}).$$

389 *Proof.* Let $x \in \mathbb{H}$. It holds that

$$\begin{aligned} 390 \quad x \in \text{argmin } f + g &\iff x = \text{prox}_{f+g}(x) \\ 391 &\iff x \in \text{Fix}(\overline{\mathcal{F}}_{f,g}(x, \cdot)) \\ 392 &\iff x = \overline{\mathcal{F}}_{f,g}(x, x) = \mathcal{F}_{f,g}(x) \\ 393 &\iff x \in \text{Fix}(\mathcal{F}_{f,g}). \end{aligned}$$

394 The proof is complete. \square

395 Let $f, g \in \Gamma_0(\mathbb{H})$ such that g is differentiable on \mathbb{H} . The classical *Forward-Backward*
 396 *algorithm* $x_{n+1} = \mathcal{F}_{f,g}(x_n)$ is a powerful tool since it provides a one-loop algorithm,
 397 only requiring the knowledge of prox_f and ∇g , that weakly converges (under some
 398 conditions on g , see [2, Section 27.3 p.405] for details) to a fixed point of $\mathcal{F}_{f,g}$, and
 399 thus to a minimizer of $f + g$.

400 From Proposition 4.1, and for all $x \in \mathbb{H}$, one can consider the one-loop algorithm
 401 (potentially weakly convergent) given by

$$402 \quad (\mathcal{A}_3) \quad \begin{cases} y_0 \in \mathbb{H}, \\ y_{k+1} = \overline{\mathcal{F}}_{f,g}(x, y_k), \end{cases}$$

403 in order to compute numerically $\text{prox}_{f+g}(x)$, with the only knowledge of prox_f and ∇g .
 404 Finally, one can also consider the two-loops algorithm

$$405 \quad (\mathcal{A}_4) \quad \begin{cases} x_0 \in \mathbb{H}, \\ x_{n+1} = y_n^*, \\ \text{where } y_n^* \text{ is given by solving the auxiliary subalgorithm} \\ \quad \begin{cases} y_{n,0} \in \mathbb{H}, \\ y_{n,k+1} = \overline{\mathcal{F}}_{f,g}(x_n, y_{n,k}), \end{cases} \end{cases}$$

406 as a potential proximal-like algorithm $x_{n+1} = \text{prox}_{f+g}(x_n)$, using only the knowledge
 407 of prox_f and ∇g .

408 Convergence proofs (under some assumptions on f and g) and numerical experiments
 409 of Algorithms (\mathcal{A}_3) and (\mathcal{A}_4) , and eventually comparisons with other known algo-
 410 rithms in the literature, should be the subject of future works.

411 **4.2. Application to sensitivity analysis for variational inequalities.** As
 412 a conclusion of the present paper, we return back to our initial motivation, namely
 413 the sensitivity analysis, with respect to a nonnegative parameter $t \geq 0$, of some

414 parameterized linear variational inequalities of second kind in a real Hilbert space H .
 415 More precisely, for all $t \geq 0$, we consider the variational inequality which consists of
 416 finding $u(t) \in K$ such that

$$417 \quad \langle u(t), z - u(t) \rangle + g(z) - g(u(t)) \geq \langle r(t), z - u(t) \rangle,$$

for all $z \in K$, where $K \subset H$ is a nonempty closed and convex set of constraints, and where $g \in \Gamma_0(H)$ and $r : \mathbb{R}^+ \rightarrow H$ are assumed to be given. Note that the above problem admits a unique solution given by

$$u(t) = \text{prox}_{f+g}(r(t)),$$

418 where $f = \iota_K$ is the indicator function of K .

419 Our aim is to provide from Theorem 2.7 a simple and compact formula for the deriva-
 420 tive $u'(0)$ under some assumptions (see Proposition 4.3 for details). Following the idea
 421 of F. Mignot in [17] (see also [13, Theorem 2 p.620]), we first introduce the following
 422 sets

$$423 \quad O_v := \{w \in H \mid \exists \lambda > 0, \text{proj}_K(v) + \lambda w \in K\} \cap [v - \text{proj}_K(v)]^\perp,$$

$$424 \quad C_v := \text{cl}\left(\{w \in H \mid \exists \lambda > 0, \text{proj}_K(v) + \lambda w \in K\}\right) \cap [v - \text{proj}_K(v)]^\perp,$$

425 for all $v \in H$, where \perp denotes the classical orthogonal of a set.

426 **PROPOSITION 4.3.** *Let $v(t) := r(t) - \nabla g(u(t))$ for all $t \in \mathbb{R}$. If the following assertions*
 427 *are satisfied:*

- 428 (i) r is differentiable at $t = 0$;
- 429 (ii) g is twice differentiable on H ;
- 430 (iii) $O_{v(0)}$ is dense in $C_{v(0)}$;
- 431 (iv) u is differentiable at $t = 0$;

then the derivative $u'(0)$ is given by

$$u'(0) = \text{prox}_{\varphi_f + \varphi_g}(r'(0)),$$

432 where $\varphi_f := \iota_{C_{v(0)}}$ and $\varphi_g(x) := \frac{1}{2} \langle D^2 g(u(0))(x), x \rangle$ for all $x \in H$.

Proof. Note that v is differentiable at $t = 0$ with

$$v'(0) = r'(0) - D^2 g(u(0))(u'(0)).$$

Note that prox_g^f is single-valued from Proposition 2.17 and Remark 2.5. From Theorem 2.7, one can easily obtain that

$$v(t) = \text{prox}_g^f(r(t)), \quad \text{and thus} \quad u(t) = \text{prox}_f \circ \text{prox}_g^f(r(t)) = \text{proj}_K(v(t)),$$

for all $t \geq 0$. Since $O_{v(0)}$ is dense in $C_{v(0)}$, we use the asymptotic development of F. Mignot [17, Theorem 2.1 p.145] and we obtain that

$$u'(0) = \text{proj}_{C_{v(0)}}(v'(0)).$$

We deduce that

$$v'(0) + D^2 g(u(0)) \circ \text{proj}_{C_{v(0)}}(v'(0)) = r'(0).$$

Since g is convex and since $C_{v(0)}$ is a nonempty closed convex subset of H , we deduce that $\varphi_f, \varphi_g \in \Gamma_0(H)$. Moreover $\partial(\varphi_f + \varphi_g) = \partial\varphi_f + \partial\varphi_g$ from Remark 2.5 and $\text{prox}_{\varphi_g}^{\varphi_f}$

is single-valued from Proposition 2.17. It also should be noted that $\nabla\varphi_g = D^2g(u(0))$. As a consequence, we have obtained that

$$v'(0) + \nabla\varphi_g \circ \text{prox}_{\varphi_f}(v'(0)) = r'(0),$$

433 that is, $v'(0) = \text{prox}_{\varphi_g}^{\varphi_f}(r'(0))$. We conclude the proof from the equality $u'(0) =$
434 $\text{prox}_{\varphi_f}(v'(0))$ and from Theorem 2.7. \square

435 *Remark 4.4.* Proposition 4.3 provides an expression of $u'(0)$ in terms of the proximity
436 operator of a sum of two closed and convex functions. Hence, it could be numerically
437 computed from Algorithm (\mathcal{A}_1), requiring the knowledge of $\text{proj}_{C_{v(0)}}$ and prox_{φ_g} .
438 Alternatively, if the convergence is proved, one can also consider Algorithm (\mathcal{A}_3)
439 requiring the knowledge of $\text{proj}_{C_{v(0)}}$ and $\nabla\varphi_g = D^2g(u(0))$.

440 *Remark 4.5.* The relaxations in special frameworks of the assumptions of Proposi-
441 tion 4.3 should be the subject of future works. In particular, it would be relevant to
442 provide sufficient conditions on K and g ensuring that u is differentiable at $t = 0$.

443 The application of Proposition 4.3 in the context of some shape optimization problems
444 with unilateral contact and friction is the subject of a forthcoming research paper
445 (work in progress).

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